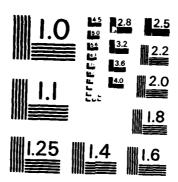
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NAVAL POSTGRADUATE SCHOOL Monterey, California



THESIS

REAL-TIME APPLICATIONS IN MULTIPROCESSOR SYSTEMS

by

M. Kadri Ozyurt

December 1983

Thesis Advisor:

Uno R. Kodres

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Real-Time Applications in Multiprocessor Systems

bу

M. Kadri Ozyurt Lieutenant J.G., Turkish Navy

Submitted in partial fulfillment of the requirements for the degree of

MASTER OF SCIENCE IN ENGINEERING SCIENCE

from the

NAVAL POSTGRADUATE SCHOOL December, 1983

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	Dean of Science and Engineering

ABSTRACT

This thesis builds a simulation model of a tactical fire control system in a real time environment, using a tightly connected multi-processing system consisting of two single board computers. The additional hardware used in this project consists of an ADM-3A video terminal with a built-in retrographics feature, an MDS microprocessor development system, an analog-to-digital convertwo sets of triplet potentiometers. The potenter, and tiometers are used to feed analog information about ownship, targetship, and gun position to the simulation model, which then evaluates and computes projected target positions and gun control parameters, and displays the results.

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I. INTRODUCTION

A. BACKGROUND

To this date, many tactical control and decision systems have been designed and implemented in various places, where the nature of the job required fast response and decision making. The NTDS (Navy Tactical Decision System), for instance, is one such system implemented for U.S. Navy ships in 1962, and is still in use today, with recent hardware modifications.

The revolutionary developments in the LSI (Large Scale Integration) and VLSI (Very Large Scale Integration) industries during the period from 1972 to 1983 have made the costs of computing much less expensive than the costs of yesterday's systems. Today's products, which are based on these innovations in large scale integration, have proved to be more reliable and more versatile than the old systems, and they can also be tailored to the needs of whatever the nature of the requirements may be. The serviceability, availability and inexpensiveness of these products, in addition to the above-mentioned features, offer both the designers and the implementers an opportunity to take advantage of this technology.

B. DISCLAIMER

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Intel corporation, Santa Clara, California: Intel, Intel 8086, iSBC 86/12A, MULTIBUS, MDS

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Micropolis Corporation, Chatsworth, California: Micropolis

Lear Siegler, Inc., Anaheim, California: ADM-3A

C. PURPOSE OF THIS THESIS

The purpose of this thesis is to create a simulation model for real time tactical systems which can be used to study the following features:

- 1. Multiprocessor system real time performance;
- 2. System reliability;
- 3. Graphics Display;
- 4. Software Engineering.

In order to carry out these objectives, certain hardware changes had to be made. The real applications necessitate the existence of an interruptdriven configuration originating from an accurate timer. The interconnections and the appropriate initializations, both in the timing (PIT) and the interrupt (PIC) circuits. were made on the iSBC 86/12A board, to give the required real time clock. After achieving the desired form of operation, the real time executive module (which synchronizes the operation of the simulation programs) was tested in an interrupt-driven environment. After testing out both the real time executive and most of the simulation system programs, the individual execution times of the simulation system programs were measured with the aid of TIMES.AID, an %INCLUDE file (see Apppendix F).

It was intended to utilize two iSBC 86/12A single board computers in order to study the real time performance of a tightly connected multiprocessing scheme. The Intel MDS (Microprocessor Development System) allows the configuration of such an expansion, through its 20 bit MULTIBUS tackplane. It was also planned to write the required software to prevent a single point failure and to gain a "graceful degradation" in the case of a malfunction in any of the single board computers.

The ADM-3A video terminal, with its built-in Retrographics unit, was utilized for graphics display purposes. The Retrographics card contains a Z-80A eight bit microprocessor. This allows the computations for high precision graphics to be done by the video terminal. That improves the efficiency by removing much of the overhead from the iSBC 86/12A single board computers. The Retrographics unit can also make drawings and erasures selectively. This improves the display time, which might be lost due to total erasures and redrawings. In other words, the selective erasing capability decreases the display time, such that the program may make partial erasures, erasures of single objects.

The programs written for the simulation model were made as modular as much as possible, to facilitate testing and maintenance, and to make room for future alterations. A procedure call was placed wherever a critical design decision was to be made. This procedure call gives the option of changing a critical design decision if one that is more efficient is designed. Structured programming and efficient data structures were meant to be utilized. Circular linked-lists are examples of such a programming technique.

D. THESIS ORGANIZATION

The thesis is organized into four chapters. The program listings developed to implement the simulation system are appended at the end of the text. The first charter covers the background, the disclaimer for the trademarks used in this thesis, the intended purpose of the project and thesis organization. The second chapter covers the system configuration and the hardware components. The third chapter deals with software modules written in both PL/I-86 and RASM-86 assembly languages. The program segments are discussed in detail in this chapter. Some information about the data structures used in the development of the software are brought up, as well as the initialization of the programmable hardware components. In the final chapter. some conclusions are presented on the work involved in the implementation of the simulation system.

II. SYSTEM HARDWARE

A. SYSTEM CONFIGURATION

The existing system hardware (see Figure 2.1) uses the Intellec microcomputer development system (MDS), which allows the expandability required to set up a multi-processor system. Within the MDS, the boards required for the operation of the system are interconnected through the MULTIBUS backplane. These boards are the following: two iSBC 86/12A boards, the front panel control board, an A to D converter board, and two interface toards for disk drives. The MDS utilizes an Intel disk drive unit which has two disk drives. Standard 8-inch IBM floppy disks are used as the removable storage media.

Fach iSBC 86/12A board had a RAM capacity of up to OFFFF hexadecimal (65535 decimal) eight bit bytes. Since the MULTIBUS is a 20-bit address bus, the address space of the whole system can be expanded up to 1 megabyte. The single board computers can address this memory space by their 20-bit address bus. Each board is so wired that the first 64 Kbyte RAM segment resides on its board (2000 through FFFFH). The 64K RAM segments on each board can be wired to be accessible from the MULTIBUS as dual ported memories in the 1-megabyte address space. One of the iSBC 86/12A boards is the master of the master-slave

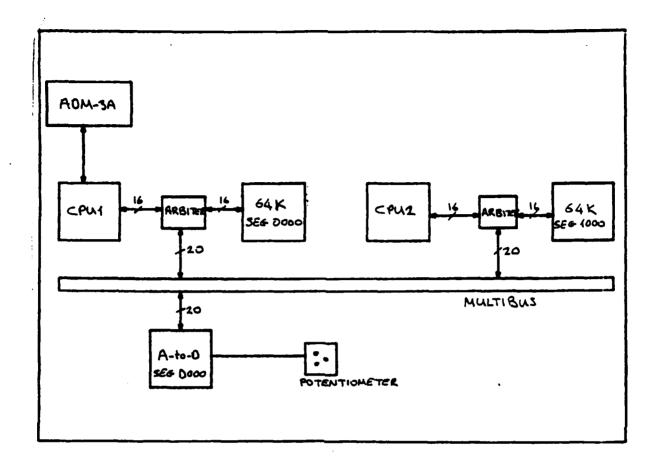


Figure 2.1 - System Interconnection

multiprocessor scheme. The master board is assigned the first 64Ksegment (00000h - 0FFFFh) where the slave board is assigned the second 64K segment (10000h-1FFFFh). In this configuration the A-to-D board occupies the segment with the addresses (D0000h-DFFFFh). The outputs of the A-tc-D converter are memory-mapped as eight bit bytes and occupy the locations DF700H through DF70FH. That is to say that they are viewed by the CPU's in the system as being

ordinary memory locations. An analogy to that is that they resemble read only memory that can only be read but not written in to.

The master iSBC 86/12A is connected to the ADM-3A video terminal through it's serial I/O connector. This is the only means of communication of the system to the outside world. The results of the simulation are fed to the video display through this interconnection. The ADM-3A video terminal has a built-in retrographics feature. consequence of this is that the ADM-3A works not only as an alphabetic terminal but also may act as a graphics device due to the fact that the retrographics card itself has a 2-80A microprocessor built in. This microprocessor allows the high precision graphics computations to be done without the need of any other external processors. In this case, ADM-3A has four operational modes, each one of which has an impact on the simulation graphics and will be discussed later in this chapter.

B. HARDWARE

In the following subsections the individual components that comprise the system hardware are presented.

1. MDS

The Intellec Microcomputer Development System (MDS) is a complete development tool which allows the integration of both microcomputer hardware and software

development. The system operates under the control of an 8086 microprocessor which supervises all system resources such as the main memory, I/O peripheral devices, and optional system facilities, such as A-to-D converters. It can support up to 7 iSBC 86/12A boards in this configuration. Some of the important boards are presented below.

a. Front Panel Control Board

The Front Panel Control Board contains circuits for controlling the front panel options. It also provides some signals for bus control, clock generation, and the bootstrap program. A bus time-out system is included to prevent the CPU from halting operation if a nonexistent memory location or an incorrect I/O port is addressed.

This board produces two types of clock pulses:

- (1) Bus Clock (10 MHz), used in Bus transactions;
- (2) Common Clock (10 MHz), used by system devices;

b. Disk Interface Boards

These two cards contain the disk controller interface for each drive in the Disk Storage Unit.

2. Single Board Computer

Intel's Single Board Computer iSBC 86/12A is used in the system. It is a member of Intel's complete line of 8- and 16-bit single board computer products and is a

complete computer system on a single printed-circuit assem-The iSBC 86/12A toard includes a 16-bit central bly. processing unit (CPU), 32K tytes (32,768 tytes) of dynamic RAM, a serial communications interface (USART), three programmable parallel I/O ports, programmable timers (PIT), priority interrupt control (PIC). Multibus interface control logic, and bus expansion drivers for interfacing with other Multibus interface-compatible expansion boards. Also included is a dual port control logic to allow the iSBC 86/12A board to act as a slave RAM device to other Multibus interface masters in the system, as is the case in this project. In the current state of the hardware, the RAM capacity of both iSBC 86/12A boards is expanded up to 64K bytes by installing an iSBC 300 Multimodule RAM option. A read only memory of 16K bytes is also added to both iSBC 86/12A boards. The important components that make up the 1SBC 86/12A board are discussed in the following subsections.

a. CPU

The iSBC 86/12A Single Board Computer is controlled by an Intel 8086 16-bit Microprocessor (CPU). The 8086 CPU includes four 16-bit general purpose registers that may also be addressed as eight 8-bit registers. In addition, the CPU contains two 16-bit pointer registers and two 16-bit index registers. Four 16-bit segment registers,

specifically: code, data, extra, and stack segment registers; allow extended addressing to a full megabyte of memory. The CPU instruction set supports many variations of addressing modes and data transfer operations, signed and unsigned 8-bit and 16-bit arithmetic including hardware multiply and divide, and logical and string operations. The CPU architecture permits dynamic code relocation, reentrant code, and instruction lookahead.

b. Serial I/O

The serial I/O port is controlled and interfaced by an Intel 8251A USART (Universal Synchronous/Asynchronous Receiver/Transmitter) chip. The USART is individually programmable for operation in most synchronous serial data transmission formats.

In the synchronous mode, the following are programmable:

- (1) Character length
- (2) Sync character (or characters)
- (3) Parity

In the asynchronous mode the following are programmable:

- (1) Character length
- (2) Baud rate factor
- (3) Stop bits
- (4) Parity

In both the synchronous and asynchronous modes, the serial I/O port features half- or full-duplex, double buffered transmit and receive capability. The USART transmit and receive clock rates are supplied by a programmable baud rate/time generator.

c. Programmable Interval Timer

Three independent, fully programmable 16-bit interval timer/event counters are provided by in Intel 8253 Programmable Interval Timer (PIT). Each counter is capable of operating in either BCD (binary coded decimal) or binary modes; two of these counters are available to the system's programmer to generate time intervals under software control. In this thesis project the counterl is used to generate timing pulses required for the real time clock operation to the system software. These pulses are sent to the PIC via the interrupt matrix as being an IR1 input request.

d. Priority Interrupt Control

The priority interrupt control (PIC) which can be programmed to respond to edge-sensitive or level-sensitive inputs, treats each true input signal condition as an interrupt request. After resolving the interrupt priority, the PIC issues a single interrupt request to the CPU. Interrupt priorities are independently programmable under

under software control. The programmable interrupt priority modes are:

- (1) Nested Priority. Each interrupt request has a fixed priority: input Ø is highest, input 7 is lowest. This mode of operation is chosen in this thesis project;
- (2) Fully Nested Priority. This is essentially the same as item (1) above, with the exception that the requesting input is not locked out and pending requests are still accepted;
- (3) Auto-Rotating Priority. Priorities are equal. The last received input becomes the lowest priority input;
- (4) Specific Priority. Software assigns the priorities;
- (5) Special Mask. Interrupt requests that are being serviced are masked out;
- (6) Poll. The CPU's internal interrupt enable is disabled. Interrupt service is achieved by a programmer-initiated Poll command.

The iSBC 86/12A board provides two sorts of interrupts which are bus vectored (BV) and non-bus vectored (NBV). The former deals with the interrupt requests from off-board sources where the latter deals with various on-board sources. The interrupt requests are fed to the

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- (1) Nested Priority. Each interrupt request has a fixed priority: input Ø is highest, input 7 is lowest. This mode of operation is chosen in this thesis project;
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- (5) Special Mask. Interrupt requests that are being serviced are masked out;
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PIC through the jumpers of the interrupt matrix, which will be discussed in the next subsection.

e. Interrupt Matrix

Interrupt requests may originate from eighteen sources without the necessity of external hardware. The interrupt matrix connects the selected source lines to a maximum of eight selected inputs of the PIC. It is an array of pins which can be connected to each other via jumper wires. There are two types of pins. The eighteen source lines constitute the input pins, where the pins that lead to the IRO through the IR7 inputs of the PIC constitute the output pins. (See Figure 2-2 for the interconnection scheme of this thesis project.)

f. Dual Port RAM

The iSBC 86/12A board has an internal bus for all on-board memory and I/O operations. Hence, local (on-board) operations do not involve the MULTIBUS interface, making it available for other iSBC 86/12A boards for a multi-processor scheme. Dual port control logic is included to interface so that the iSBC 86/12A board can function as a slave RAM device (or common memory) when not in control of the Multibus interface. The CPU has priority when accessing on-board RAM. After the CPU completes its read or write operation, the controlling bus master is is allowed to access RAM and complete its operation. Where both the CPU and the controlling bus master have the need

to write or read several bytes or words to or from the onboard RAM, their operations are interleaved. CFU access, the on-board RAM addresses are assigned from the bottom up of the 1-megabyte address space; i.e., 00000-ØEFFFh. The slave RAM address decoding logic includes jumpers and switches to allow positioning the on-board RAM into any 64-K segment of the 1-megabyte system address The slave RAM can be configured to allow either 16K, 32K, 48K, or 64K access by another bus master, with the installation of the iSBC 300 Multimodule RAM. In this thesis project all of the 64K-byte memory of the slave iSBC 86/12A is made accessible to the master. Furthermore. both iSBC 86/12A boards are configured to occupy the first 128K section of the 1-megabyte address space.

3. A-to-D Converter Board

This board is electrically and mechanically compatible with any iSBC 86/12A board and with MDS. Both the anolog input and output systems are contained on a single printed circuit board that is treated as ordinary memory locations by the CPU (memory mapping). This board simply gets the analog signals form the potentiometers and converts them to the digital signals compatible with TTL standards. The output of the A-to-D converter is one byte per potentiometer input, which varies from +127 to -128. This is the maximum value range a fixed binary (7) variable can assume in the PL/I language, by definition. So, the

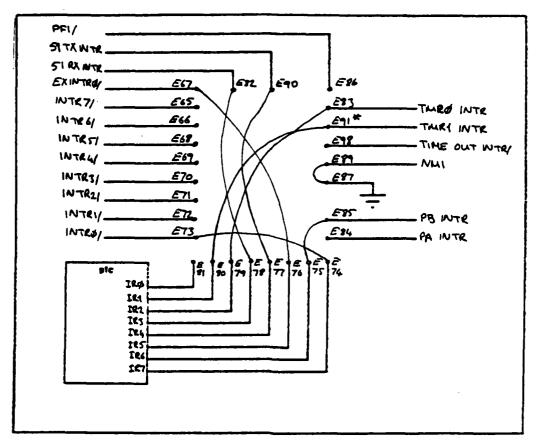


Figure 2.2 -Interrupt Matrix Interconnection

software system thinks of the potentiometer value as changing between +127 and -128. The A-to-D converter board is configured to occupy the segment D000: (14th 64K portion of the 1-megabyte address space).

4. ADM-3A Terminal

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This terminal is the only means by which the operator communicates with the system. It is connected to the system with the master iSBC 86/12A board's serial I/O connector. It is an interactive device which is used to enter, display, and send information to a host computer.

and to receive and display information from that computer. The information exchange between the terminal and the computer is made at different baud rates, ranging up to 19200. In this scheme, a 9600 baud rate is used. The keyboard contains 59 keys. The display memory is a RAM which is capable of holding 1920 characters. Data characters are displayed on 24 or 12 equally-spaced rows, each consisting of 80 columns.

5. RG-512 Retrographics Card

The RG-512 Retrographics printed circuit board is added to the ADM-3A terminal to extend the data and graphics display capabilities with the aid of a Z-80A built-in microprocessor. The RG-512 employs the bit map method of storing graphic images. This information is in a digital memory as a rectangular array of bits. Each bit is mapped onto the CRT screen and can cause a bright spot to be displayed. The RG-512 displays graphs and pictures by writing the proper bits into graphics memory. One of the important features of the RG-512 is the ability to erase portions of the screen selectively. This is desirable when the application requires the use of dynamic displays employing motion or rotation to convey information. The RG-512 has four modes of operation. These are the ADM-3A Alpha Mode, the Alpha Mode, the Point Mode, and the Vector Mode. The first one is equivalent to the operation of ADM-3A without RG-

512. The latter three modes make use of the bit map method.

6. Potentiometers

Two sets of triplet potentiometers are used as simulating analog sensor information sources. They feed anolog signals to the input of the A-to-D converter board, varying between -5V and +5V.

III. IMPLEMENTATION OF THE SOFTWARE

A. GENERAL INFORMATION

1. Modularity

A modular and extensible simulation program is aimed at simplifying the debugging and testing phase and at facilitating possible alterations. The hierarchy of the modularity is composed of a head module and four secondlevel modules connected to the head module. These secondlevel modules are the initialization module, the simulation system module, the real time executive module. and the dynamic debugging tool module. These modules separately compiled PL/1-86 and RASM-86 programs. main module is further subdivided into third-level modules to gain a finer granularity of modularity. In the programs, two useful special features of PL/I-86 are Those are %INCLUDE and %REPLACE statements. used. Ву those statements, global declarations that are the same in the scope of the simulation program need not be declared within each and every module. Instead, they are grouped together in the GLOBALS. INP declaration file. The %REPLACE statement allows constants to be declared as in the other high-level languages like Pascal, such that the value of the constants can be changed without having to go through every program segment in which they occur.

2. Data Structures

Linear arrays and arrays of structures (records) are used in the simulation program as data structures. These data structures are then linked to each other to establish circular linked lists. Figure 3.1 (on the next page) explains the general picture of the circular linked lists in the simulation program. Fixed size data structures with fixed binary pointers are used in the program, rather than pointer data, to avoid the dynamic system overhead and to retain the benefit of random access capabilities inherent to linear arrays.

The structure SHIP has two pointers, in addition to the fields that hold specifications about the ships in the area. Those pointers, PTR and LINK-SHIP have different purposes for different ships. SHIP (1), for instance, being the ownship points to two different circular linked lists. PTR points to the enemy ship's circular linked list by pointing to the target ship which is engaged (tracked) for the sea battle. The other pointer, LINK-SHIP points to the friendly ships circular list. The PTR field of other ships, on the other hand, points to another circular linked list, WAKE, to record their past positions that will be used for tracking and display purposes, where LINK-SIP points to the other ships in their category. The reason for using circular lists in this program is the ease with

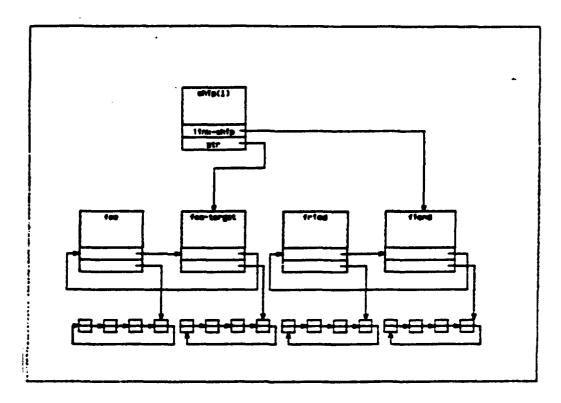


Figure 3.1 The structure of the circular linked lists

which one traverses through the lists and which does not necessitate the use of another external pointer.

3. Other Features

The PL/I exception handlers (ON body statements) are used extensively in the dynamic debugging module, and in various interactive parts of the program, to intercept the error conditions that might be raised during the testing and execution of the program. The ON body statements are contained in the ERRHAND.AID file, which is an %INCLUDE file. Upon receiving the control through a raised error condition, the statements in the file prompt

the user and give the control over to the REENTRY.PLI interactive debugging tool.

One of the exceptional features of the simulation program is the use of non-local goto statements which are unacceptable in structured style of programming. It is an inevitable requirement, by the PL/I language, to suppress the raised error conditions by a non-local goto statement. It is also used in some parts of the dynamic debugging module, in order to by-pass the flow of control over to the debugging program when the optional ERRORON boolean switch is closed. Explicit comments are offered wherever non-local goto statements are used, to avoid confusing the reader.

B. SOFTWARE FUNCTIONAL DESCRIPTION

In the following sections, the structure of the modules and the programs that belong to those modules are described. The program listings are presented as appendices.

1. Head Module

a. WAR.PLI

This main procedure is the head node of the hierarchical structure of the procedures used to modularize and structure the implementation of the simulation program. It contains two call statements, one of which is to the initialization module, to set up the tactical database

and to initialize various external variables that are used throughout the simulation program. The other call is for passing the control to the real time executive module where the control stays for the rest of the program execution. A listing of WAR.PLI is presented as Appendix A.

2. Initialization Module

a. INITVARS.PLI

This PL/I routine, when called from the main procedure, constructs the tactical database in interactive manner. It first initializes the pool of available SHIP and WAKE modes for later use. It then gets interval of time which is used to update the information about ships in the tactical area and other time dependent functions. This time interval must be equal to the period of the real-time interrupts which depend on the timing constants used during the initialization of PIT. Detailed information will be presented later under the hardware initialization section. Then the control proceeds to establish the tactical database interactively. In this session, the initial information about azimuth, range, friend or enemy are written to the proper fields. Then, a circular linked list of four nodes is composed and the PTR field is made to point to that list. Finally, dependent on whether friend or foe, that particular node is added to the appropriate circular linked list. INITVARS.PLI makes use of various internal subroutines for linked list

operations. Those subroutines simply extract a node from the pool of available nodes. After establishing the linked lists, the INITVARS.PLI initializes external variables that are used throughout the program. A listing of INITVARS.PLI is presented in Appendix B.

3. Simulation System Module

This module is composed of four PL/I programs which perform the simulation under the control of the real-time executive module. The following subsections describe the functional description of those programs. The listings of the program segments that comprise this module are presented in Appendix C.

a. TACTICAL.PLI

This routine has the highest priority among the system module programs. It first updates the position of each ship in the tactical area by calculating the relative velocity and multiplying that with the time interval, which is the period of timing interrupts that occur every 250 milliseconds. The control then proceeds to calculate the future positions of the ships, for those ships which have been in the area for more than 4 seconds, and which are included in the enemy ship circular linked list. The routine uses the polynomial least squares curve fitting method with Legendre Polynomials. The coefficients are pre-calulated for the position of the ship one second after the time of calculations, based on the past four wake

points. TACTICAL.PLI also calculates the trajectory of the travelling projectile if the gun is fired.

b. DISPLAY.PLI

This routine, which is invoked every second, is the interface of the simulation program for the ADM-3A screen. It simply traverses the circular linked lists and generates the appropriate display objects for the ships, their aim points which the own ship aims at, and the gun aim point. It also displays the travelling projectile if it is fired by the system. The routines for generating the objects are internal for the DISPLAY.PLI. Another routine, TRANSLATE, translates the cartesian coordinates to the stream of characters that represent the x and y grid coordinates of the RG-512 Retrographic Screen Memory. Finally, the internal routine DRAW puts the generated objects on the screen in vector mode. According to the key variable. D. it either sets the data level to white and puts the object on the screen, or sets the data level to black and makes selective erasures.

c. STATUS.PLI

This routine, which is invoked at every second, is the interface of the system to the ADM-3A video terminal keyboard. The commands for the system are read from the keyboard by calling the serial I/O chip (USART) interface assembly program. KEYBOARD, which will be presented under Miscellaneous Assembly Routines.

The are four boolean variables used in the simulation program. Those are ENGAGED, MAGNIFIED. FIRED. ERRORON. The STATUS.PLI sets these variables and according to the value of the parameter passed to the assembly interface routine. ENGAGED (with "E") shows if the system is engaged to any target (for TACTICAL.FLI). MAGNIFIED (with "M"), which is used in DISPLAY.PLI, shows which display scale is being used and what the reference point of the display is. Usually, the display on the screen is relative to own ship. But that can be changed so that the ship engaged is at the center to the screen by setting MAGNIFIED true. FIRED (with "F") boolean variable is used to commence the ballistic calculations and display. It is used both in TACTICAL.PLI and DISPLAY.PLI. Finally, ERRORON (with "D") is used to transfer the control to the dynamic debugging module.

d. IDLE.PLI

This program is the idle routine for the system, as it waits for a 250 millisecond timing interrupt to occur. It is the interface program for the six potentiometers which are used as sensors for the own ship velocity vector, the known ship velocity vector, and the gun elevation and bearing. The velocity vectors are composed of speed and course components. The IDLE.PLI gets this information by calling the Analog-to-Digital Converter interface assembly program, ATOD. The control then

proceeds to convert this information, which is in the fixed binary (7) form in range (-128, +127), to appropriate coordinate values, e.g. 50 knots maximum speed and the true azimuth between 0 and 360 degrees. This routine calculates own ship velocity components in cartesian coordinates for later use by TACTICAL.PLI. It also makes the initial ballistic computations for the gun.

4. Real Time Executive Model

This module works as the interrupt handler for the real time interrupts that are initiated by programmable hardware components every 250 milliseconds. It is invoked by the WAR.PLI the first time and interrupts thereafter. It responds to the timing interrupts, which tell the system that data must be collected at this point in time. module then resolves the priorities of the simulation system module programs and arbitrates the flow of control during the execution of the system. The real time module makes use of the operating executive primitives, which are presented in the following subsections. The labels P1 through P4 are associated with the simulation system modules TACTICAL.PLI, DISPLAY.PLI, STATUS.PLI, and IDLE.PLI, respectively. The listing of the programs included in the Real Time Module are presented in Appendix D.

a. ARBITER.A86

This assembly language program is the real "workhorse" of the entire system. It first allocates stack areas for four simulation system during the assembly time. Upon invocation by the main procedure, it initializes the programmable hardware components and transfers the control to the P4.PLI process, which in turn calls IDLE.PLI repetitively until the first timing interrupt occurs. interrupt entry point PROCO, where the process switching starts, is entered by the interrupt software. At this point of execution. the external variable fourtheve, which signals the system that 250 millisecond event has occured, is updated to the new value by incrementing it by one. After storing the state of the program which is interrupted during its execution, ARBITER.A86 invokes SCHEDULE.PLI to obtain the name of the ready program that has the highest priority. If there are none. the interrupted program is resumed. If there is any ready process of higher priority, then ARBITER.A86 loads the process state and gives the control over to it. During this process switching, the upper boundaries of the stack areas are checked for a possible stack overflow, which could happen if the time interval was not sufficiently large for the system module routines to finish execution before the next interrupt comes. ARBITER.A86 has a second entry point, STORESTATUS, for the synchronization primitive AWAIT.PLI to enter when the correct number of interrupts for the calling synchronization primitive have not yet occured.

b. AWAIT. PLI

This synchronization primitive is invoked as as an operating system primitive, by any process, P1 through P3. AWAIT compares the value of the external variable FOURTHEVC to the threshold value of the calling process to see if it is greater or equal to the value at which the process is to proceed. If not, then it calls the STORESTATUS entry of ARBITER.A86 to relinquish the control to the awaiting ready process, or to the P4.PLI that calls the IDLE.PLI simulation system program as the system idling routine.

c. SCHEDULE.PLI

This synchronization primitive is called by ARBITER.A86 to return the name of the highest priority ready process. It does that simply by identifying the first ready process on the list. Because the scheduler scans the list in the descending priority order, the highest priority ready process will automatically be scheduled.

d. THRESH.PLI

This routine, when invoked by processes P1 through P3, increments the corresponding thresholds in an external one dimensional array called THRESHOLD. This table is used by the AWAIT.PLI and SCHEDULE.PLI

synchronization primitives to decide whether or not a process is ready for execution.

e. P1.PLI

This process is basically an infinite loop. Within this loop, there are three subroutine calls. The routine first makes a call to AWAIT.PLI to see if it is time for it to proceed. If not, the control doesn't come back again; instead, the current state of the process is stored by ARBITER.A86 and the highest priority ready process is executed. If it is the time, the control proceeds to call the simulation system module program TACTICAL.PLI. After that, a call to THRESH.PLI is made, where the threshold value that is allocated to P1.PLI is incremented by the proper value. When the infinite do loop repeats itself, the call to the AWAIT.PLI will indicate that the process TACTICAL.PLI is not yet ready for execution and control is transfered to the highest priority ready process.

f. P2.PLI

This process is identical to P1.PLI in form except the call is to DISPLAY.PLI instead of to TACTICAL.PLI.

g. P3.PLI

This process is identical to P1.PLI in code except the call is to STATUS.PLI.

h. P4.PLI

This process is similar to P1 through P3 in the structure described above. There is only one call in the infinite do loop, which is to IDLE.PLI. This routine is always ready for execution and, basically, repeats itself until the next timing interrupt comes along.

5. Miscellaneous Assembly Routines

There are few machine dependent functions that cannot be accomplished by the high-level language PL/I-86. Assembly routines were written to interface the PLI-86 programs with the hardware of the 8086 microprocessor. These assembly routines are included within the main body ARBITER.A86. There are two parameter passing of conventions from PL/I-86 to the assembly language routines. In the first one, there is only one argument passed in the accumulator. as in a function call. In the subroutine calls. which is the case here. the address of the VECTOR that contains the pointers to the actual parameters is passed in the BX register. The following subsections give some descriptions about those assembly routines. Appendix E shows the listings of the modules.

a. KEYBOARD.A86

This routine is invoked by STATUS.PLI to read the keyboard. It is written so that the keyboard status is read to see if a key had been pressed instead of waiting indefinitely until a key was depressed, as would be the

first reads in the status of the serial I/O interface chip (USART) to see if a character has been received from the keyboard. If it has, then the character is read and placed into a corresponding character variable, which is the formal parameter in the subroutive invocation. If there is no character received since the last attempt to read, the ASCII equivalent of 0 is put in key. The reason for that is that the character 0 is not being used as a command.

b. ATOD.A86

This assembly routine is called by IDLE.PLI to read the first six Analog-to-Digital Converter Board outputs. The reason to write this assembly routine is that the Analog-to-Digital Converter ports are memory mapped to be in the segment D000H. The PL/I function UNSPEC works for those memory locations which are included within the first 64K bytes of memory. The assembly routine sets the proper segment and source index registers to point to those locations and makes an ordinary read operation. This value is then put in the formal parameter passed to the PLI-86 routine.

c. RINGBELL.A86

This assembly routine simply sends a bell character to the video terminal and causes a bell sound to ring. This is equivalent to sending a control G in PL/I-86.

The only difference is that it can be used in other assembly routines.

d. WAIT.A86

This assembly routine reads in the status of the I/O interface chip and waits until the transmitter buffer is empty; i.e. the character which had been in the buffer is received by the video terminal and an acknowledgement signal is sent back to the interface chip. This routine is used by assembly routines that put out a message.

e. SUSPEND.A86

This routine simply resets the interrupt bit of the program status word (PSW) to disable the 8086 CPU from acknowledging the interrupts. It is used by the dynamic debugging system to stop the real time clock.

f. RESUME.A86

This routine first sets the interrupt flag to enable the 8086 CPU to respond to the interrupt requests. It also resets the PIT that is used to generate the timing clock pulses for the PIC.

D. INITIALIZATION OF THE PROGRAMMABLE COMPONENTS

The iSBC 86/12A board has three programmable hardware components, which were described in Chapter II. In the following subsections, the initialization sequences for

those hardware components which produce the real time synchronization are described.

1. USART Programming

The 8251A USART converts parallel output data into virtually any serial output data format. The USART also converts serial input data into the parallel data format. Prior to starting to transmit or to receive data, the USART must be loaded with a set of control words. These control words, which define the complete functional operation of the USART, must immediately follow a reset (internal or external). There are two types of control words, namely, a Mode instruction and a Command instruction. Once the Mode instruction has been sent, the Command instruction can be sent at any time prior to a read/write operation. The following assembly code is used to initialize the USART read/write mode:

MOV AL,37H OUT ØØDAH.AL

During the course of execution, the serial I/O interface routine executes the following assembly code to read the USART status and to read the receiver buffer if any character has been received:

IN AL,00DAH AND AL,02H JZ KEYBOARD1 IN AL.00D8H

2. PIT Programming

The 8253 PIT has three independent counters. They are Counter 0, Counter 1, And Counter 2. The input clock frequency is 22.1184 MHz supplied by a crystal oscillator. The input clock frequencies can be adjusted via jumpers. In this project, Counter 1, with the default factory jumper connection (E59-E60), is used. The imput frequency is 153.6 KHz. It is chosen to work in mode 0 so that it gives an interrupt on terminal count. The formula:

N = TC

where

N is the count value for the counter

T is the desired interrupt time interval in seconds

C is the internal frequency (Hz)

From the above formula, the count number for the counter 1 is found to be 38400 decimal (9600 hex). Since it is not possible to express this number in four decimal digits, the binary count mode is selected. The counter is initialized by sending a mode control word, followed by a down-count number. Once the counter is initialized, sending the down-count number resets it to the start condition. The following sequence of code is used to program the PIT:

VOM AL,50H OUT ØØD6H,AL MOV AL.ØØH ØØD2H.AL OUT MO V AL,60H ØØD6H,AL OUT VOM AL,96H OUT ØØD2H.AL

3. PIC Programming

The 8259A PIC is programmed in the nested mode. The master PIC with no slaves is accepted. For this particular situation, the initialization words 1,2, and 4 are sent. The initialization word 2 is set to represent the interrupt vector address. This is the address that the control is given when the interrupt 1 occurs (Ø4H). An interrupt mask byte is used to mask out the irrelevant interrupts for the purpose of this thesis. The PIC can be reset after each interrupt simply by sending the EOI (end of interrupt) status byte to the appropriate address. The initialization sequence is as follows:

CLI
MOV AL.13H
OUT ØØCØH,AL
MOV AL.2ØH
OUT ØØCZH
STI
MOV AL.ØFDH
OUT ØØCZH,AL

To reset the PIC.

MOV AL, 20H OUT 00C0, AL

E. ASSEMBLY. COMPILING AND LINKING

The assembly language code was written in RASM-86 and assembled by using the RASM-86 Assembler. This assembler produces relocatable files that can then be linked with other separately compiled or assembled object files by Link 86. This linker accepts three types of input files. Those are the object file, litrary file, and/or an input file. Input files are very useful tools in that they include input command lines by an input file instead of writing the command line each time the programs are to be linked. The PL/I-86 compiler is used to compile PL/I programs. This compiler requires a 128 Kbyte RAM, as opposed to the PL/I-80 compiler, which requires 48K RAM.

F. TESTING

The hierarchical simulation program modules were designed and tested in a top-down manner. An extensive dynamic debugging module was used in the testing of individual modules (see Appendix F). The testing phase is first started with writing a skeletal model for the real time executive module. The PL/I-86 output statements (stubs), which printed some appropriate numbers, were inserted in the places where the simulation system module programs were invoked. The DDT86 (Dynamic Debugging Tool) was used to test and debug this skeleton program. Since it was a real-time interrupt driven program, the interrupt

enable bit (I) of the program status word was reset to zero (0) to control the flow of the program. Some error checkings are inserted in the main assembly language program, ARBITER.A86, to see whether the process switching was done correctly, or to see if any of the stacks allocated to the synchronization primitives have overflowed.

After testing the real-time executive model, the second phase was the testing of the subroutines that are used by the simulation system module programs. Appendix G shows an example for testing a subroutine.

The third phase of the testing was to test and debug the simulation system program. The dynamic debugging module was used for this purpose. This module is composed of PL/I-86 %INCLUDE files and external PL/I-86 programs. The %INCLUDE files are inserted into the various parts of the program being tested. The code of this debugging system is bordered with comment lines from the main body of the program it tests. It is not visible to the program, i.e., it brings its local and global variables with the LOCALS.AID declaration %INCLUDE file. It is possible to manipulate the system's external variables through the debugging system module, IDLE.PLI.

The PL/I-80 had been used in the early stages of the testing phases because of the existence of redundant Intel 8080 based systems in the Naval Postgraduate Micro-Lab.

But in the later stages, it was realized that this had caused some problems, due to some incompatibilities between the PL/I-80 and PL/I-86 systems. These deficiencies of the PL/I-86 have necessitated the testing and debugging of the programs to be done in the PL/I-86 based systems. The need for such a thorough debugging system module was then realized.

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IV. CONCLUSIONS

original objectives of the thesis have been The large extent. accomplished. to a The hardware interconnections intended to promote real time operations have been successful. A test model, which comprised the real time executive module and the PL/I put statements, instead of the simulation system modules, was developed to check the system's operation and timing. Correct results appeared on the screen. Then, the testing and debugging of the simulation system programs showed that their algorithms and operations were correct, with the exception of the simulation system program, DISPLAY.PLI. The testing of DISPLAY.PLI showed that the objects to be displayed on the video screen were not successfully put on the screen. One error, which was an automatic conversion error, was found in the routine TRANSLATE and corrected by using a step variable. However, an error still exists in the body of the procedure DRAW.

Since the testing phase of the DISPLAY.PLI has not been accomplished, the objective of synchronizing two iSBC 86/12A boards could not be accomplished. The intended purpose for this objective was to make the second iSBC 86/12A jump to a waiting loop with the initial power start up interrupt (reset), to load the assigned simulation program segment to the common RAM, and to direct it to the

beginning of that program segment, through a bus vectored interrupt.

A hierarchical and modular program model was constructed through the use of data structures and structured programming. The modularization allows possible future changes to the programs.

The simulation system program constructed in this thesis could be a basis for further enhancements towards a complete fire control system, or a related tactical simulation system, due to its modularized nature.

The dynamic debugging module designed to test and debug the simulation system modules can be used for the purpose of testing any other programs, simply by changing the names and formats of the variables that the debugging system manipulates. The advantage of using the %REPLACE and %INCLUDE pre-processor statements, which are peculiar to the PL/I-86 version of the subset G, makes such an implementation feasible.

APPENDIX A

HEAD MODULE PROGRAM LISTINGS

A. WAR.PLI /* Prog Name WAR.PLI December 83 Date Written by M. Kadri Ozyurt For Thesis Advisor Professor Kodres This is the main procedure of the Purpose It invokes the initialization modular simulation program. module to set up the target database and to initialize the external variables used throughout the simulation program. */ WAR: PROCEDURE OPTIONS (MAIN); /*external procedures*/ DCL (INITVARS, ARBITER) ENTRY; /*this call to the initialization module initializes the simulation system*/ CALL INITVARS; /*this call gives the control over to the real time executive module*/ CALL ARBITER;

END WAR;

B. GLOBALS.INP

: GLOBALS.INP Prog Name : December 83 Date : M. Kadri Ozyurt Written by : Thesis For Advisor : Professor Kodres : This %include file contains the declara-Purpose tions of the global variables used in the simulation program DCL (COUNTER.SECONDS.MINUTES.HOURS.WAKE PTR. SHIP PTR, AVAILSHIP, AVAILWAKE, P, Q, NUMBERSHIPS, NODE, TARGET, KNOWN) FIXED BIN(7) EXTERNAL, FOURTHEVC FIXED BIN(15) EXTERNAL , D FIXED BIN(7) EXTERNAL, (DT.T_PRIME.T_OF,T) FLOAT EXTERNAL, (I,J,XX,YY) FIXED BIN(15). CURRENTPROC FIXED BIN(7) EXTERNAL . (ENGAGED.MAGNIFIED.FIRED.ERRORON.DONE) BIT(1) EXTERNAL . KEY CHARACTER (1) EXTERNAL THRESHOLD (0:2) FIXED BIN (15) EXTERNAL . ARG(0:5) FIXED BIN(7) EXTERNAL, (VX OWN, VY OWN, VX TARGET, VY TARGET, VX REL, VY REL.VX ROUND. VY ROUND, VR) FIXED DECIMAL EXTERNAL, ALPHA FIXED DECIMAL EXTERNAL . (AX SUM.BX SUM.CX SUM.AY SUM.BY SUM,CY SUM, AX, BX, CX, AY, BY, CY, X_AT5, Y_AT5, R, DX_DT AT5. DY DT AT5, DR DT AT5, X OFFSET, Y OFFSET, M) FIXED DECIMAL EXTERNAL. (01,02)(5) FIXED DECIMAL EXTERNAL, 1 SHIP (MAX SHIPS) EXTERNAL. 2 VELOCITY. 3 COURSE FIXED DECIMAL INIT (0.0). 3 SPEED FIXED DECIMAL INIT (0.0). 2 POSITION. 3 AZIMUTH FIXED DECIMAL INIT(0.0), 3 RANGE FIXED DECIMAL INIT (0.0).

```
2 COORDINATES.
        3 X FIXED DECIMAL
                             INIT (0.0).
        3 Y FIXED DECIMAL
                             INIT (0.0).
    2 AIM,
        3 X AIM FIXED DECIMAL
                                 INIT(\emptyset.\emptyset).
        3 Y AIM FIXED DECIMAL
                                INIT(Ø.Ø).
    2 COUNT FIXED BIN (7) INIT(0),
    2 NUMBER FIXED BIN (7) INIT(0),
    2 PTR FIXED BIN (7) INIT(0).
    2 LINK SHIP FIXED BIN (7) INIT(0),
    2 FRIEND BIT(1) INIT(FALSE).
1 OBJECT (MAX SHIPS) EXTERNAL.
    2 LOCATIONS.
        3 U (0:10) FIXED EIN(15) INIT((11) -1),
        3 V (0:10) FIXED BIN(15) INIT((11) -1),
    2 AIMS.
        3 U AIM (0:10) FIXED BIN(15) INIT((11) -1),
        3 V AIM (0:10) FIXED BIN(15) INIT((11) -1),
    2 GUN,
        3 U GUN (0:10) FIXED BIN(15) INIT((11) -1),
        3 V GUN (0:10) FIXED BIN(15) INIT((11) -1),
    2 WAKES,
        3 U_WAKE (0:10) FIXED BIN(15) INIT((11) -1);
        3 UU_WAKE (0:10) FIXED BIN(15) INIT((11) -1),
        3 V WAKE (0:10) FIXED BIN(15) INIT((11) -1),
        3 VV WAKE (0:10) FIXED BIN(15) INIT((11) -1).
1 GUN EXTERNAL.
    2 POSITION.
        3 AZ FIXED DECIMAL
                             INIT(0.0).
        3 ALT FIXED DECIMAL
                             INIT(0.0).
    2 COORDINATES,
        3 X GUN FIXED DECIMAL
                                 INIT(0.0),
        3 Y GUN FIXED DECIMAL
                                 INIT(0.0).
1 WAKE(4) EXTERNAL.
    2 COORDINATES.
        3 X WAKE FIXED DECIMAL
                                  INIT(\emptyset.\emptyset),
        3 Y WAKE FIXED DECIMAL
                                  INIT(0.0),
    2 LINK WAKE FIXED BIN(7) INIT(0).
(SUSPEND, RESUME, ARBITER, INITYARS) ENTRY,
KEYBOARD ENTRY (CHARACTER(1)),
ATOD ENTRY (FIXED BIN(?).FIXED BIN(?));
```

C. CONST.INP

Prog Name : CONST.INP

Date : December 83

For : Thesis

Advisor : Professor Kodres

Purpose : This %include file contains the constant declarations used throughout the program.

*/

```
TREPLACE
                          /*max number of wake nodes*/
/*max number of ship nodes*/
MAX_WAKE BY 4,
MAX SHIPS BY 2,
                          /*max number of variables*/
MAXVARS BY 82,
OWN BY 1,
                          /*ownship indicator*/
RMAX BY 25000.0,
                          /*max gun range*/
MAISQ BY 1.073E+09.
                          /*max argument for SQRT */
                          /*max number for fixed (15)*/
TOP BY 32767.0,
                          /*gravitational con.yd/sec2*/
G BY 10.7246,
VM BY 518.0.
                          /*muzzle velocity*/
                          /*x coord. for center*/
A BY 512,
                          /*y coord. for center*/
B BY 390,
K BY 1.40625,
L BY 4.513,
                          /*azimuth proportionality c.*/
                          /*speed proportionality c.*/
                          /*definition for 360 degree*/
TWO PI BY 360.0,
                          /*definition of pi rad/180deg*/
PI BY 3.1416,
 00 BY 1.0,
                          /*legendre poly. of zero deg.*/
TRUE BY '1'B.
                          /*boolean true*/
FALSE BY '0'B,
                          /*boolean false*/
/*in linked list terminology*/
VECMOD BY TO
                          /*
                          /* "
                          /* "
 ALPHA4010 BY
                          /* "
 ALPHAJA BY
WHITEMOD BY
                          /*
 BLACKMOD BY
                          /*
```

APPENDIX B

INITIALIZATION MODULE PROGRAM LISTINGS

```
INITVARS.PLI
                  : INITVARS.PLI
Prog Name
                  : December 83
Date
Written by
                  : M. Kadri Ozyurt
                  : Thesis
For
                  : Professor Kodres
Advisor
                  : This routine prompts the user to give
Purpose
the time interval (dt) and constructs the tactical circular
linked list in an interactive manner. The control then
proceeds to initialize the external variables.
initvars:proc external;
  dcl
          %include'const.inp';
          %include'globals.inp';
  /*this iterative loop initializes the pool of available
wake nodes*/
  do i=1 to max wake-1;
          link_wake(1)=1+1;
  end;
  link wake(max wake)=nil;
/*this sequence initializes the pool of available ship
nodes*/
  do i=1 to max ships-1;
          link_ship(i)=i+1;
  end;
  link_ship(max_ships)=mil;
  put skip list('Enter the time interval (dt) in seconds');
```

```
/* the following block is an interceptor for a too large
input value*/
 on overflow begin;
put list('*** too large, try again');
          goto init1;
 end;
init1:
 put skip list('>');
 get list(dt);
 revert overflow;
 put skip list('Construction of the tactical database');
 numberships=0;
 done=false;
/* this procedure call gets a ship node from the pool of
available ships and assigns its address to ship pointer*/
  ship ptr=getship();
 do while ( done);
          put skip list('Enter the position of ', numberships,
                   th ship in true azimuth and in yards');
/*the following
                  two on condition bodies are for the input
line at init2*/
          on error begin;
                  put list('*** bad value.try again');
                  goto init2;
          end;
          on fixedoverflow tegin;
                  put list('*** too large, try again');
                  goto init2;
          end;
init2:
          put skip list('>');
          get list(azimuth(ship_ptr),range(ship_ptr));
          revert error;
          revert fixedoverflow;
          put skip list('Friend or foe (F/E)?');
          put skip list('>');
          get list(key);
/* the following sequence adds the ship node to the
appropriate circular linked list, friend ships or enemy
ships, according to the friend boolean value entered*/
          if (key='f')!(key='f') then do;
               friend(ship_ptr)=true;
              if link_ship(own)=nil then
                  link_ship(ship_ptr)=ship_ptr;
              else
                  link_ship(ship_ptr)=link_ship(own);
              /*end if*/
```

```
link ship(own)=ship ptr;
               end /*do*/;
           else do;
                friend(ship_ptr)=false;
                if ptr(own)=nil then
                    link_ship(ship_ptr)=ship_ptr;
               else
                    link_ship(ship_ptr)=ptr(own);
                /*end i?*/
                ptr(own)=ship ptr;
           end /*if*/;
/* the following procedure call and the iterative loop get
available nodes from the pool, construct a circular linked list of four nodes, and assign the address of the list to
      ptr pointer of the ship node which
                                                      is
constructed*/
           wake_ptr=getwake();
           q=wake_ptr;
           do i=1 to 4;
                    p=wake_ptr;
                    wake_ptr=getwake();
                    link_wake(p)=wake_ptr;
           end /*do*/;
           link_wake(wake_ptr)=q;
           ptr(ship_ptr)=wake_ptr;
           number(ship_ptr)=numberships;
           put skip list ('Would you like to enter another
                              'sbir (Y/N)?');
           put skip list('>');
           get list(key);
if (key='Y')!(key='y') then do;
                    ship_ptr=getship();
                    if ship_ptr=nil then
                             done=true;
                    /*end if*/
                    end:
           else
                    done=true;
           /*end if*/
  end /*do*/;
  seconds=0;
  minutes=0;
  hours=0;
  wake ptr=0;
  target=0;
  t_prime=0.0;
  t of=0.0;
  t=0.0;
```

```
link ship(own)=ship_ptr;
end /*do*/;
          else do;
               friend(ship_ptr)=false;
               if ptr(own)=nil then
                  link ship(ship_ptr)=ship_ptr;
              else
                  link_ship(ship_ptr)=ptr(own);
               /*end i ?*/
               ptr(own) = ship ptr;
          end /*if*/;
/* the following procedure call and the iterative loop get
available nodes from the pool, construct a circular linked
list of four nodes, and assign the address of the list to
          pointer of the ship node which
     ptr
constructed*/
          wake_ptr=getwake();
          q=wake_ptr;
          do i=1 to 4;
                  p=wake_ptr;
                  wake_ptr=getwake();
                  link wake(p)=wake ptr;
          end /*do*/;
          link_wake(wake_ptr)=q;
          ptr(ship_ptr)=wake_ptr;
          number(ship_ptr)=numberships;
          put skip list ('Would you like to enter another',
                            ship (Y/N)?');
          put skip list('>');
          get list(key);
          if (key='Y')!(key='y') then do;
                  ship_ptr=getship();
                  if ship_ptr=nil then
                           done=true;
                   /*end if*/
                  end:
          else
                  done=true;
          /*end 1f*/
  end /*do*/;
  seconds=0;
  minutes=0;
  hours=0;
  wake_ptr=0;
  target=0;
  t prime=0.0;
  t of=0.0;
  t=0.0;
```

```
friend(own)=true;
  known=2;
                           /* engaged ship no by the sensor */
  r_offset=0.0;
  y_offset=0.0;
  m=50.0:
  fourthevc=0;
  1=0;
  j=0;
  currentproc=4;
  engaged=false;
 magnified=false;
  fired=false;
  erroron=false;
 key='0';
  threshold(\emptyset)=1; threshold(1)=4; threshold(2)=4;
  VI OWn=0.0;
  vy_own=0.0;
  vr target=0.0;
  vy_target=0.0;
  vr rel=0.0;
  vy_rel=0.0;
  vx round=0.0;
  vy_round=0.0;
  vr=0.0;
  alpha=0.0;
  ax_sum=0.0; bx_sum=0.0; cx_sum=0.0;
  ay sum=0.0; by sum=0.0; cy sum=0.0;
  ar=0.0;tr=0.0;cr=0.0;
  ay=0.0;by=0.0;cy=0.0;
  x_at5=0.0;y_at5=0.0;
  r=0.0;
  dx_dt_at5=0.0; dy_dt_at5=0.0; dr_dt_at5=0.0; o1(1)=1.0; o1(2)=0.5; o1(3)=0.0; o1(4)=-0.5; o1(5)=-1.0;
  02(1)=1.0; 02(2)=-0.5; 02(3)=-1.0; 02(4)=-0.5; 02(5)=1.0;
/*getship. when invoked. extracts a node from the pool of
available and returns a pointer value pointing to that
node. It puts an error message if there is no available
node*/
getship:procedure returns (fixed bin(7));
  /* dcl */
          %include 'globals.inp';
  if availship=nil then do;
          put skip list('No more available ship nodes');
          return(nil);
          end /*do*/;
  else do;
          node=availship;
          availship=link_ship(availship);
  end /*if*/;
```

```
return(node);
end getship;
/*getwake does the same function as getship except for the
operations made are on wake nodes*/
getwake:procedure returns (fixed bin(7));
   /* dcl */
            %include 'globals.inp';
   if availwake=nil then do;
    put skip list('No more available wake nodes');
            return (Ø);
            end /*do*/;
   else do;
            node=availwake;
            availwake=link_wake(availwake);
   end /*if*/;
   return(node);
end getwake;
end initvars;
```

APPENDIX C

A. TACTICAL.PLI

SIMULATION SYSTEM MODULE PROGRAM LISTINGS

```
Prog Name
                     : TACTICAL.PLI
Date
                     : December 83
Written by
                     : M. Kadri Ozyurt
For
                     : Thesis
Advisor
                     : Professor Kodres
                     : This external routine calculates and
Purpose
updates the positions of the ships in the tactical area and the future positions of the ships that belong to
enemy ships circular linked list pointed to by ptr(own),
and calculates the trajectory of the travelling projectile
if fired
*/
TACTICAL: PROCEDURE EXTERNAL;
/*
  DCL
    %INCLUDE 'CONST.INP';
    %INCLUDE 'GLOBALS.INP';
/*following sequence of code updates the present positions
of the ships in the tactical area*/
  DO I=2 TO NUMBERSHIPS;
       VI_TARGET = SPEED(I) * SIND (COURSE(I));
VI_TARGET = SPEED(I) * COSD (COURSE(I));
      VX_REL = VX_TARGET - VX_OWN;
VY_REL = VY_TARGET - VY_OWN;
      WAKE PTR = PTR(I);
```

```
X WAKE(WAKE PTR) = X(I);
       Y WAKE (WAKE PTR) = Y(I);
X(I) = X(I) + VX_REL * DT;
       Y(I) = Y(I) + VY REL * DT;
                                         /*ptr(i) points to*/
       PTR(I) = LINK WAKE(WAKE PTR);
       IF COUNT(I) < 4 THEN
                                           /* the oldest wake*/
           COUNT(I)=COUNT(I)+1;
       /*END IF*/
  END /*DO*/:
/*calculating the future positions (aim
                                                   points)
here by using the least squares method with legendre
polynomials. The coefficients are pre-calculated according
to the fifth second including the zeroth second*/
FILTERING:
  TARGET = PTR(OWN);
  DONE=TRUE;
       (TARGET=NIL) THEN
           DONE=FALSE:
  /*END IF*/
  DO WHILE ( DONE);
       IF COUNT(TARGET)=4 THEN
         BEGIN;
           AX SUM = X(TARGET) * 00;
                                           /*Leg. poly. 0 deg.*/
           BX_SUM = X(TARGET) * O1(5);/*
                                                          1
           CX_{SUM} = X(TARGET) * 02(5);/*
                                                                 */
           AY_SUM = Y(TARGET) * 00;
           BY SUM = Y(TARGET) * O1(5);
           CY = SUM = Y(TARGET) * 02(5);
           WAKE_PTR = PTR(TARGET);
           J = \overline{1};
           DO WHILE (LINK WAKE(WAKE PTR) = PTR(TARGET));
                AX_SUM = A\overline{X}_SUM + X_WAKE(WAKE_PTR) * 00;
                BX_SUM = BX_SUM + X_WAKE(WAKE_PTR) * O1(J);
                CX_SUM = CX_SUM + X_WAKE(WAKE_PTR) * O2(J);
                AY SUM = AY SUM + Y WAKE (WAKE PTR) * OO;
BY SUM = BY SUM + Y WAKE (WAKE PTR) * O1(J);
CY SUM = CY SUM + Y WAKE (WAKE PTR) * O2(J);
                WAKE PTR=LINK WAKE (WAKE PTR);
                J = \bar{J} + 1;
           END /*DO*/;
           AX = AX_SUM / 5.;
           BI = 2.0 + BI SUM / 5.;
CI = 2.0 + CI SUM / 7.;
           AY = AY_SUM / 5.;
           BY = 2.0 * BY SUM / 5.;
           CY = 2.0 * CY SUM / 7.;
```

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```
X AT5 = AX - 1.5*BX + 3.5*CX;
           Y = AY - 1.5*BY + 3.5*CY;
INCONVENIANCE:
       begin block is inserted to avoid the complications
which might arise from the automatic conversions*/
        BEGIN;
             DCL
                   (XSTEP.YSTEP) FIXED.
                   (XF.YF.RSQD) FLOAT;
              IF (ABS(X AT5)>TOP)! (ABS(Y AT5)>TOP) THEN
                   RSQD=-1.0;
               ELSE DO;
                   XSTEP=BINARY(X_AT5);
YSTEP=BINARY(Y_AT5);
                   XF=FLOAT(XSTEP);
                   YF=FLOAT (YSTEP);
                    RSOD=XF*XF+YF*YF;
               END /*IF*/;
              IF (RSQD<0.0)! (RSQD>MAXSQ) THEN
                   R=0.0;
              ELSE
                   R=SQRT(RSQD);
               /*END IF*/
           END INCONVENIANCE;
           IF (R=0.0) ! (R>RMAX) THEN DO;
                    X_AIM(TARGET)=0.0;
                Y AIM(TARGET)=0.0;
END /*DO*/;
           ELSE DO:
                    ALPHA = ASIN(G*R/VM**2) / 2.0;/*IN RADS*/
                    VR = VM + COS(ALPHA);
                    T_PRIME = R / VR;
DX_DT_AT5 = 3.0*CX - 0.5*BX;
                    DY DT AT5 = 3.0*CY - 0.5*BY;
                    DRDTAT5 = (XAT5 * DXDTAT5 +
                   Y_AT5 * DY_DT_AT5) / R;
T_OF = (R + DR_DT_AT5 * T_PRIME) / VR;
                   X = AIM(TARGET) = X = AT5 + DX DT AT5 * T OF;
                    Y_AIM(TARGET) = Y_AT5 + DY_DT_AT5 * T_OF;
           END /*IF*7;
      END /*IF*/;
      TARGET=LINK SHIP(TARGET);
                                           /*next target?*/
      IF TARGET=PTR(OWN) THEN
           DONE=TRUE;
       /*END IF*/
  END /*DO*/;
```

```
ROUNDTRACK:

/*the ballistic calculations start here*/

IF (FIRED) THEN

BEGIN;

VX_REL = VX_ROUND - VX_OWN;

VY_REL = VY_ROUND - VY_OWN;

I_GUN = X_GUN + VX_REL * DT;

Y_GUN = Y_GUN + VY_REL * DT;

T = T - DT;

IF T<=0 THEN DO;

PUT LIST('G');

FIRED = FALSE;

END /*IF*/;

END /*IF*/;
```

END TACTICAL;

```
B. DISPLAY.PLI
 /#
Prog
      Name
                   : DISPLAY.PLI
Date
                   : December 83
Written by
                   : M. Kadri Ozyurt
For
                   : Thesis
Advisor
                   : Professor Kodres
Purpose
                       This routine first puts the time in
hours, minutes, and seconds. The control then proceeds to
call DRAW subroutine in an iterative loop to erase the old
objects. It then calculates the positions of the objects
relative to either ownship
                                or the ship that has been
targeted according
                      to MAGNIFIED. Then it calls
                                                     DRAW
display the objects.
DISPLAY: PROCEDURE EXTERNAL;
/*DCL*/
       %INCLUDE 'CONST.INP';
       %INCLUDE 'GLOBALS.INP';
  PUT LIST('^]^ 'X');
                            /*ENTER ALPHA MODE*/
  IF MINUTES = 0 THEN
       BEGIN:
           PUT LIST ('^[= %');
           PUT EDIT (HOURS)(F(2));
HOURS = HOURS+1;
           IF HOURS=24 THEN
              HOURS = 0;
           /*END IF*/
  END /*IF*/:
  IF SECONDS=0 THEN
       BEGIN;
           PUT LIST ('^[= (');
           PUT EDIT (MINUTES) (F(2));
           MINUTES = MINUTES+1;
           IF MINUTES=60 THEN
              MINUTES = 0;
           /*END IF*/
  END /*IF*/;
  PUT LIST ( '^[= +');
  PUT EDIT (SECONDS)(F(2));
```

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```
SECONDS = SECONDS+1:
  IF SECONDS=60 THEN
      SECONDS = \emptyset;
  /*END IF*/
  PUT LIST( 'M'X');
                                   /*HOME CURSOR*/
/*the following calls erase the objects from the screen*/
  DO I=1 TO NUMBERSHIPS;
          CALL DRAW(OBJECT(I).U.OBJECT(I).V.D);
          CALL DRAW(OBJECT(I).UU WAKE.OBJECT(I).VV WAKE.D);
          CALL DRAW(OBJECT(I).U AIM.OBJECT(I).V AIM.D);
  END /*DO*/;
  CALL DRAW(OBJECT(1).U GUN, OBJECT(1).V GUN.D);
       following sequence converts the coordinates
objects to the grid coordinates of the screen and generates
the sequences of the coordinates for the objects*/
 DO I=1 TO NUMBERSHIPS;
          XX=A+BINARY((X(I)-X_OFFSET)/M);
YY=B+BINARY((Y(I)-Y_OFFSET)/M);
          IF (FRIEND(I)) THEN
             CALL GENFRIEND(XX.YY.OBJECT(I).U.OBJECT(I).V);
          ELSE
             CALL GENFOE(XX.YY.OBJECT(I).U.OBJECT(I).V);
          /*END IF*/
          CALL GENWAKE(XX.YY,OBJECT(I).U WAKE,OBJECT(I).V WAKE);
          XX=A+BINARY((X_WAKE(PTR(I))-X OFFSET)/M);
          YY=B+BINARY((Y_WAKE(PTR(I))-Y_OFFSET)/M);
          IF COUNT(I)=4 THEN
             CALL GENWAKE(XX.YY.OBJECT(I).UU WAKE.
                                          OBJECT(I).VV WAKE);
          XX=A+BINARY((X_AIM(I)-X OFFSET)/M);
          YY=B+BINARY((Y AIM(I)-Y OFFSET)/M);
               ((XX=A)&(YT=B)) THEN
          IF
                IF I=1 THEN
                   CALL GENOURAIM(XX.YY.OBJECT(I).U AIM.
                                             OBJECT(I).V AIM);
               ELSE
                   CALL GENAIM(XX.YY.OBJECT(I).U AIM.
                                            OBJECT(I).V_AIM);
                /*END IF*/
           /*END IF*/
  END /*DO*/;
  XX=A+BINARY((X GUN-X OFFSET)/M);
  YY=B+BINARY((Y_GUN-Y_OFFSET)/M);
  CALL GENGUN(XX,YY,OBJECT(1).U_GUN,OBJECT(1).V_GUN);
```

```
/*the following sequence draws the objects by calling
routine DRAW*/
  D=1;
  DO I=1 TO NUMBERSHIPS;
      CALL DRAW(OBJECT(I).U.OBJECT(I).V.D);
      CALL DRAW(OBJECT(I).U_WAKE,OBJECT(I).V_WAKE,D);
      IF (ENGAGED) THEN CALL DRAW(OBJECT(I).U_AIM,
                                            OBJECT(I).V_AIM,D);
  END /*DO*/;
  IF FIRED THEN CALL DRAW(OBJECT(1).U GUN,OBJECT(1).V GUN,D);
/* the following procedures produce the sequence of screen
grid coordinates for various objects*/
GENFRIEND:PROC(X.Y.U.V);
  DCL
           (X.Y) FIXED BIN(15),
           (U.V)(0:10) FIXED PIN(15);
  U(Ø)=X;
                   V(\emptyset) = Y - 8;
  U(1)=X-6;
                   V(1)=Y-3;
  U(2)=X-6;
                   V(2)=Y+3;
  U(3)=X;
                   V(3) = Y + 8;
                   V(4)=Y+3;
  U(4) = X + 6;
  U(5) = X + 6;
                   V(5)=Y-3;
  U(6)=X;
                   V(6)=Y-8;
  U(7)=-1;
                   V(7)=-1;
END GENFRIEND;
GENFOE: PROC(X,Y,U,V);
  DCL
           (X,Y) FIXED BIN(15),
           (U.V)(0:10) FIXED BIN(15);
                    V(Ø)=Y-4;
  U(\emptyset) = X + 8;
  U(1)=X-8;
                   V(1)=Y-4;
  U(2)=X;
                   V(2)=Y+8;
  U(3) = X + 8;
                   V(3)=Y-4;
  U(4)=-1;
                   V(4)=-1;
END GENFOE;
```

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```
GENWAKE:PROC(X.Y.U.V);
  DCL
            (X,Y) FIXED BIN(15),
(U,V)(0:10) FIXED BIN 15);
                      V(Ø)=Y;
V(1)=Y;
  U(0)=X;
U(1)=X;
                      V(2)=-1;
  U(2)=-1;
END GENWAKE:
GENOURAIM: PROC(X,Y,U,V);
  DCL
             (X,Y) FIXED BIN(15),
             (U.V)(Ø:10) FIXED EIN(15);
                       V(\emptyset)=Y;
   U(Ø)=X+8;
                       V(1)=Y;
   U(1)=X-8;
                       V(2)=Y;
   U(2)=X;
  U(3)=X;
U(4)=X;
U(5)=-1;
                       V(3)=Y-8;
                       V(4)=Y+8;
V(5)=-1;
END GENOURAIM;
GENAIM:PROC(X,Y,U,V);
   DCL
             (X,Y) FIXED BIN(15),
             (U.V) (0:10) FIXED PIN(15);
                       V(Ø)=Y+4;
   U(3)=X+4;
   U(1)=X-4;
                       V(1)=Y-4;
                       V(2)=Y;
   U(2)=X;
                       V(3)=Y-4;
   U(3)=X+4;
U(4)=X-4;
                       V(4)=Y+4;
                       V(5)=-1;
   U(5)=-1;
 END GENAIM;
```

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```
GENGUN:PROC(X,Y,U,V);
  DCI.
           (X.Y) FIXED BIN(15),
           (U.V)(0:10) FIXED PIN(15);
                    V(Ø)=Y-1;
  U(2) = X + 1;
  U(1)=X+1;
                    V(1)=Y+1;
  U(2)=X-1;
                    V(2)=Y+1;
  U(3) = X - 1;
                    V(3)=Y-1;
  U(4) = X + 1;
                    V(4)=Y-1;
  U(5) = -1;
                    V(5)=-1:
END GENGUN;
/*this procedure receives two arrays and a key variable as
paramaters, and either displays the object or erases it#/
DRAW:PROC(U.V.D) EXTERNAL;
  DCL
           (U,V)(0:10) FIXED BIN(15), (I,J,H,D) FIXED BIN(?),
           RUB CHAR(1) EXTERNAL.
           C(7) CHAR(1),
           Z1(0:3) CHAR(1) BASED(P).
           Z(0:3) BIT(8),
           P POINTER:
  P=ADDR(Z);
                    /*Z and Z1 share same location hereon*/
  I=0;
  DO WHILE (I<11);
           IF (D=1)
                /*enter vector set level wnite*/
THEN PUT LIST( ^ ] [a');
                /*enter vector set level black*/
ELSE PUT LIST(']^[',RUB);
           /*END IF*/
           DO H=1 TO 5;
                    IF (U(I)<\emptyset) THEN DO;
                              PUT LIST( 'M X [=
                                                    ′);
                              RETURN;
                    END /*IF*/:
             /*this call translates the coordinates to the
                stream of bits*/
                    CALL TRANSLATE(U(1),V(1),Z);
                    I=I+1;
```

```
/*this put statement puts out the bit streams
                as being characters*/
                   PUT EDIT((Z1(J) DO J=0 TO 3)) (4A(1));
          END /*DO*/:
        /*the following two statements get the status of
        the screen. The status is not sent back until the
        screen is ready.*/
          /*HANDSHAKE*/
  END /*DO*/;
 PUT LIST( "M X ");
                            /*BACK TO ALPHA*/
TRANSLATE: PROCEDURE (X,Y,Z);
  DECLARE (X.Y) FIXED BIN(15),
          Z(0:3) BIT(8),
          T FIXED BIN(15).
          S1 BIT (16),
           SS BIT (7).
          S BIT (8),
           I FIXED BIN(7);
  I = DIVIDE(Y.32.8);
  SS=BIT(I,7);
  S='0'B || SS;
Z(0) = '00100000'B ! S;
  T = Y - (Y/32) * 32;
  S1 = BIT (T,16);
S = SUBSTR(S1,8,8);
  Z(1) = '01100000'B ! S;
  I = DIVIDE (X.32.8);
 SS=BIT(I,7);
S='0'B || SS;
Z(2) = '00100000'B ! S;
T = I - (I/32) * 32;
  S1 = BIT(T,16);
  S = SUBSTR (S1,8,8);
  Z(3) = '010000000'B ! S;
END TRANSLATE;
END DRAW;
```

END DISPLAY;

```
STATUS . PLI
                   : STATUS.PLI
Prog
      Name
Date
                   : December 83
                   : M. Kadri Ozyurt
Written
         bу
For
                   : Thesis
                   : Professor Kodres
Advisor
Purpose
                   : this routine calls the assembly routine
KEYBOARD to read the keyboard to set the boolean variables
used in other routines
STATUS: PROCEDURE EXTERNAL;
/*
 DCL
     %INCLUDE 'CONST.INP';
               'GLOBALS.INP';
     ZINCLUDE
      CALL KEYBOARD (KEY);
      IF (KEY='Q')!(KEY='q') THEN
          STOP;
      ELSE IF (KEY='E')!(KEY='e') THEN
          ENGAGED = TRUE;
      ELSE IF ENGAGED & ((RANK(KEY)>48)&(RANK(KEY)<=57))
        THEN BEGIN:
          I=RANK(KEY)-48;
          IF NUMBER(I)=0 THEN
                  LINK_SHIP(I)=I;
          ELSF IF FRIEND(I) THEN DO;
                  CALL REMOVENODE(LINK SHIP(OWN),I);
                  CALL ADDNODE(PTR(OWN),I);
                  FRIEND(I)=FALSE;
                  END;
          PTR(OWN) = I;
          END;
      ELSE IF (KEY='R') ! (KEY='r') THEN
          ENGAGED=FALSE;
      ELSE IF (KEY='M')!(KEY='m') THEN DO;
```

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```
/*set the scale to 1/200*/
           MAGNIFIED=TRUE:
           M=200.0;
           IF ENGAGED THEN DO;
           /*set reference as the target*/
                    X OFFSET=X(PTR(OWN));
                    Y OFFSET=Y(PTR(OWN));
                    END /*DO*/:
           ELSE DO;
           /*ownship is the reference*/
                    X OFFSET=0.0;
                    Y^{-}OFFSET=\emptyset.\emptyset;
                    END /*DO*/;
           /*END IF*/
      END;
      ELSE IF (KEY='T')!(KEY='t') THEN DO;
      /*set the scale back to normal (1/50)*/
                    MAGNIFIED=FALSE:
                    M=50.0;
                    X_OFFSET=0.0;
Y_OFFSET=0.0;
         END /*DO*/;
      ELSE IF (KEY='F')!(KEY='f') THEN
           FIRED = TRUE;
      ELSE IF (KEY='D')!(KEY='d') THEN
           ERRORON=TRUE:
      ELSE IF (KEY='S')!(KEY='S') THEN
           SIGNAL ERROR(1);
      /*END IF*/
      KEY='0';
/*this routine removes the node pointed by QQ from
                                                               the
circular linked linked list pointed by PP*/
REMOVENODE: PROC (PP.QQ);
  DCL
           (PP,QQ) FIXED BIN(7);
           ZINCLUDE GLOBALS. INP ';
  P=PP;
  P=LINK SHIP(P);
DO WHILE ( (LINK SHIP(P)=QQ));
P=LINK SHIP(P);
  END;
  LINK SHIP(P)=LINK SHIP(QQ);
END;
```

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D. IDLE.PLI
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```
/*
                 : IDLE.PLI
Prog
     Name
                 : December 83
Date
Written by
                 : M. Kadri Ozyurt
For
                 : Thesis
                 : Professor Kodres
Advisor
Purpose
                 : This routine reads the A/D converter
output to get the velocity vectors of ownship and
selected ship, and the gun information as azimuth and
elevation. It then converts this information to real world
values. It calculates ownship speed which will be used to
find relative speeds later. It then computes the maximum
       cartesian coordinates of the splash point, and time
of flight corresponding to the current gun position
IDLE: PROCEDURE EXTERNAL;
/*
 DCL
     %INCLUDE 'CONST.INP';
     %INCLUDE 'GLOBALS.INP';
     DO D=Ø TO 5;
         CALL ATOD (D, ARG(D));
     END /*DO*/;
     this point the A/D output values are fixed bin(7)
        The following sequence converts those to
decimal values*/
     COURSE(OWN) = ARG(Ø);
     SPEED(OWN)=ARG(4);
     COURSE(KNOWN)=ARG(2);
     SPEED(KNOWN)=ARG(3);
     AZ=ARG(1);
     ALT=ARG(5);
/*the following sequence converts A/D values to real
```

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```
COURSE(KNOWN) = COURSE(KNOWN) * K;
      AZ = AZ * K;
      IF COURSE(OWN)<0.0 THEN
          COURSE(OWN) = COURSE(OWN) + TWO PI;
      IF COURSE(KNOWN)<0.0 THEN
          COURSE(KNOWN) = COURSE(KNOWN) + TWO PI;
      IF AZ<0.0 THEN
          AZ = AZ + TWO PI;
       IF ALT>90.0 THEN
          ALT = 90.0;
          SPEED(OWN) = SPEED(OWN)/L;
          SPEED(KNOWN) = SPEED KNOWN) / L;
/*ownship speed computations*/
      VX \cap WN = SPEED(OWN) * SIND(COURSE(OWN));
      VY OWN = SPEED(OWN) * COSD(COURSE(OWN));
/*when not have fired, the following makes the ballistic
computations*/
      IF
           FIRED THEN
         BEGIN;
          T OF = 2.0 * VM * SIND(ALT) / G;
          VR = VM + COSD(ALT);
          R = VR * T OF;
X_AIM(OWN) = R * SIND(AZ);
          TAIM(OWN) = R * COSD(AZ);
          X GUN = 0.0;
          Y GUN = \emptyset.0;
          V\bar{X}_ROUND = VR + SIND(AZ);
          VY = ROUND = VR + COSD(AZ);
          T = T OF;
      END /*IF*/:
END IDLE:
```

APPENDIX D

REAL TIME EXECUTIVE MODULE LISTINGS

A. ARBITER.A86

;Prog Name : ARBITER.A86 ;Date : December 83 ;Written by : M. Kadri Ozyurt

; For : Thesis

;Advisor : Professor Kodres

; Purpose : This program contains all the assembly ; routines used by the simulation system. It initializes all ; programmable hardware components, responds to the timing ; interrupts, and increment the FOURTHEVC used throughout the ; simulation model program. Upon receiving interrupt ; requests , it performs process switching by storing the ; state of interrupted process in the stack area allocated ; for the processes and by restoring the highest ready ; process given by SCHEDULER

;GLOBALS

A CANADA MANAGA MANAGA

DGROUP GROUP FOURTHEVC, CURRENTPROC FOURTHEVC DSEG COMMON FOURTHEVC1 DW Ø CURRENTPROC DSEG COMMON CURRENTPROC1 DB 4

CSEG EXTRN SCHEDULE: FAR EXTRN P1:FAR EXTRN P2:FAR EXTRN P3:FAR P4:FAR EXTRN PUBLIC ARBITER PUBLIC STORESTATUS RINGBELL PUBLIC PUBLIC KEYBOARD PUBLIC ATOD SUSPEND PUBLIC PUBLIC RESUME

; EQUATES

```
EOU
  INT1
                   84H
                                              ; INTR1 JUMP ADDR.
  INT3
           EQU
                   8CH
                                              ; INTR3
                                                       JUMP ADDRESS
  PIC1
           EQU
                   ØCØH
                                              ;PIC COMMAND OUTPUT PORT1
  PIC2
           EQU
                   ØC2H
                                              FPIC COMMAND OUTPUT PORTS
  ICW1
           EQU
                   13H
                                              ;PIC COMMAND WORD1
  ICW2
         · EQU
                   2ØH
                                              ;PIC COMMAND WORD2
  ICW4
           EQU
                   ØDH
                                              ;PIC COMMAND WORD4
  MASK1
           EQU
                   ØFDH
                                              FIC MASK BYTE
           EQU
  EOI
                   2ØH
                                              :END-OF-INTERRUPT BYTE
  CNTR1
           EQU
                   50H
                                              PIT MODE CONTROL BYTE
  CNTR2
           EQU
                   6ØH
  PORTC
           EQU
                   ØD6H
                                              ;PIT CONTROL PORT
           EQU
  COUNT
                   ØD2H
                                              ;PIT COUNT # OUTPUT PORT
  CNTRLO EOU
                   ØØH
                                              ;PIT COUNT # LO BYTE
  CNTRHI EQU
                   96H
                                              ;PIT COUNT # HI BYTE
  READWR
          EQU
                   37H
  RXRDY
          EQU
                   Ø2H
                                              ;USART STATUS MASK(READ)
  TXRDY
          EQU
                   Ø1E
                                              ;USART STATUS MASK(WRITE)
  PORTIO EQU
                   ØD8H
                                              JUSART I/O PORT
  PORTST
          EQU
                   ØDAH
                                              JUSART STATUS PORT
  SECCONV EQU
                   ØDØØØH
                                              ;A/D CONVTR PORT SEGMENT
  OFFCONV EQU
                   ØF700H
                                              ;A/D CONVTR PORT OFFSET
  STACKSIZE EQU
                   100H
          EQU
  LF
                   ØAH
  CR
           EQU
                   ØDH
  BEL
           EQU
                   07 H
  FS
                   1CH
           EOU
  ESC
           EQU
                   1BH
 FF
          EQU
                   ØCH
           EQU
                   188
  CAN
  PALSE
          EQU
  TRUE
           EQU
                   NOT FALSE
ARBITER:
          PUSH
                   DS
           CLI
                                              ;DISABLE INTR'S
           MOV
                   AX.Ø
           MOV
                   DS,AX
                                              ;SET SEGREG TO Ø
           MOV
                   BX, INT1
           MOV
                   WORD PTR [BX], OFFSET PROCØ ; INT1 JMP ADDRESS
           INC
           INC
                   BI
          MOV
                   WORD PTR [EX],CS
          POP
                   DS
           MOV
                   BX.OFFSET STACKTBL+STACKSIZE-2
          VOM
                   CS:[BX].CS
```

BX, STACKSIZE

ADD

```
CS:[BX],CS
           MOV
                    BX, STACKSIZE
CS:[BX], CS
           ADD
           MOV
                    BX, STACKSIZE
           ADD
                    CS:[BX],CS
           MOV
           CLI
                                                 :INIT. PIC TO
                    AL, ICW1
           MOV
                                                 ; EDGE-TRIG .. SINGLE PIC
                    PIC1.AL
           OUT
                    AL, ICW2
           MOV
                                                 ; INT1 ADDR IS Ø4H
                    PIC2,AL
           OUT
                    AL.ICW4
           MOV
                                                 ; NOT F. NESTED, NORM. EOI
                    PIC2.AL
           OUT
           STI
           MOV
                    AL, MASK1
                                       ;ONLY INT1 IS ALLOWED
                    PIC2,AL
           OUT
                                                 ;INIT. PIT
           MOV
                    AL, CNTR1
                                                 ; SELECT MODE Ø , CNTR 1 ; CLK FREQ. IS 153.6 KHZ
           OUT
                    PORTC .AL
                                                 COUNT-DOWN VALUE 9600H
           MOV
                    AL, CNTRLO
                                                 : WHICH GIVES AN INTR AT
                    COUNT,AL
           OUT
           MOV
                    AL, CNTR2
                    PORTC, AL
           OUT
                                                 : EVERY FORTH OF A SEC.
                     AL. CNTRHI
           VOM
                     COUNT, AL
           OUT
                     AL, READWR
           MOV
                     PORTST.AL
           OUT
           MOV
                     AX,CS
                                                SET STACK SEG. TO CODE
                     SS,AX
           MOV
                     BP,3*STACKSIZE
           MOV
                     SP, STACKTBL [BP]
           MOV
           JMP
                     P4
PROCØ:
                     AX
           PUSH
                     PX
           PUSH
                     CI
           PUSH
                     AL. CNTR2
           VOM
                                                 RESET COUNTER
                     PORTC, AL
           OUT
                                                 RESET THE CNTR.
           MOV
                     AL, CNTRHI
           OUT
                     COUNT .AL
                                                 RESET PIC
           MOV
                     AL, EOI
                     PIC1.AL
            OUT
                                                  ;SP->INTERRUPTED IP
            ADD
                     SP,6
                                                                      IP
                     BX
                                                  :BX->
            POP
                                                  ; AX->
                                                                       CS
                     AX
            POP
                                                 ;SP->INTERRUPTED;PUSH
                                                                     IP
            POP
                     CI
                                                                       CS
                     AX
            PUSH
                                                                       IP
                                                  ; PUSH
                     BX
            PUSH
                     CX
            PUSH
```

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```
SUB
                    SP.6
                                               ;SP->PUSHED CX
           POP
                    CX
           POP
                    BX
                                               RESTORE BX
                    AX
           POP
                                               ; RESTORE AX
           INC
                    FOURTHEVC1
           JMP
                    STORESTATUS1
STORESTATUS:
                    SP
           DEC
                    SP
           DEC
                    AX
           PUSH
                    SP,4
           ADD
           POP
                    AX
                                               ;AX->INTERRUPTED IP
           PUSH
                    CS
                    AX
           PUSH
           DEC
                    SP
           DEC
                    SP
           POP
                    AX
                                               FRESTORE AX
           PUSHF
           CLI
STORESTATUS1:
           PUSH AX ! PUSH BX ! PUSH CX ! PUSH DX
           PUSH BP ! PUSH SI ! PUSH DI ! PUSH ES
           CALL
                    SCHEDULE
           VOM
                    NEWPROC, AL
           MOV
                    AL, CURRENTPROC1
           CMP
                    AL.1
                    OUT1
           JNZ
           MOV
                    BP.1
                    OUT4
           JMP
OUT1:
                    AL.2
           CMP
           JNZ
                    OUT2
           MOV
                    BP.STACKSIZE+1
           JMP
                    OUT4
OUT2:
           CMP
                    AL.3
           JNZ
                    OUT3
           VOM
                    BP,2*STACKSIZE+1
                    OUT4
           JMP
OUT3:
           CMP
                    AL,4
           JNZ
                    OUT5
           MOV
                    BP.3*STACKSIZE+1
           JMP
                    OUT 4
OUT4:
                    BP, OFFSET STACKTBL
           ADD
           CMP
                    SP, BP
                    OUT6
           JNA
                    BP
           DEC
                     [BP],SP
           MOV
           JMP
                    LOADPROC
OUT5:
                    AL,30H
           ADD
           MOV
                    OUTSAL, AL
           MOV
                    DX.OFFSET OUT5MESS
           JMP
                    ERRORMESS
```

```
AL,30H
OUT6:
           ADD
           MOV
                    OUTGAL, AL
           VOM
                    DX.OFFSET OUTSMESS
                    ERRORMESS
           JMP
LOADPROC:
                    AL, NEWPROC
           MOV
           CMP
                    AL,1
           JNZ
                    LOUT1
           MOV
                    BP,Ø
                    RETURNPT
           JMP
LOUT1:
                    AL,2
           CMP
                    LOUTS
           JNZ
                    BP, STACKSIZE
           MOV
           JMP
                    RETURNPT
LOUT2:
           CMP
                    AL,3
                    LOUT3
           JNZ
                    BP,2*STACKSIZE
           MOV
                    RETURNPT
           JMP
LOUT3:
           CMP
                    AL,4
                    LOUT4
           JNZ
           MOV
                    BP,3*STACKSIZE
           JMP
                    RETURNPT
LOUT4:
           ADD
                    AL,3ØH
                    LOUTAAL, AL
DX. OFFSET LOUTAMESS
           MOV
           MOV
           JMP
                    ERRORMESS
RETURNPT:
           MOV
                    CURRENTPROC1, AL
           MOV
                    SP.STACKTBL [BP]
           POP
                    ES : POP DI ! POP SI ! POP BP
                    DX ! POP CX ! POP BX ! POP AX
           POP
           POPF
           STI
           RETF
                                                        OUT
;THIS ROUTINE
                             SYSYTEM CALL
                                                  PUT
                                                             ERROR
                  MAKES A
                                             TO
; MESSAGES
ERRORMESS:
                    CL,9
           MOV
           INT
                    224
           MOV
                    CL.Ø
           MOV
                    DL,1
           INT
                    224
           RET
```

THE STACK AREAS AND VARIABLE DEFINITIONS

```
STACKTBL:
                   OFFSET STACKTBL+STACKSIZE-22
           DW
                    STACKSIZE-8
           RS
           DW
                   FALSE
                   OFFSET P1
           DW
           RS
                   OFFSET STACKTBL+2*STACKSIZE-22
           DW
           RS
                    STACKSIZE-8
           DW
                    FALSE
           DW
                    OFFSET P2
           RS
                   OFFSET STACKTBL+3*STACKSIZE-22
           DW
           RS
                    STACKSIZE-8
           DW
                    FALSE
                   OFFSET P3
           DW
           RS
           DW
                   OFFSET STACKTBL+4*STACKSIZE-22
           RS
                    STACKSIZE-8
           DW
                   FALSE
           DW
                   OFFSET P4
           RS
NEWPROC
           DB
OUT 5MESS
           DB
                    FS.ESC.FF.CAN
           DB
                    'RETURN FROM OUT5. AN UNKNOWN CURRENT PROCEDURE:'
OUT 5AL
           DB
                    FS, ESC, FF, CAN
OUT6MESS
           DB
           DB
                    'RETURN FROM OUT6. STACK OVERFLOW FOR THE PROC .: '
                   Ø, '$', Ø
FS, ESC, FF, CAN
OUTSAL
           DB
LOUT4MESS DB
                    'RETURN FROM LOUT4. AN UNKNOWN NEW FROCEDURE:'
           DB
                   0,'$',0
LOUT4AL
           DB
```

END

B. AWAIT.PLI

THE PROPERTY OF CANADAS AND SOUTH AN

```
: AWAIT.PLI
Prog Name
                   : December 83
Date
        рÀ
Written
                   : M. Kadri Ozyurt
                   : Thesis
For
Advisor
                   : Professor Kodres
                      This synchronization primitive checks
Purpose
     threshold value for the calling process by
the
                                                  comparing
                              value with FOURTHEVC
the corresponding threshold
returns the control either to the calling process.
                                                     if its
threshold value is equal to and greater than FOURTHEVC, or
else transfers the control to ARBITER.A86
```

C. SCHEDULE.PLI

```
: SCHEDULE.PLI
Prog Name
                   : December 83
: M. Kadri Ozyurt
Date
Written by
For
                   : Thesis
                   : Professor Kodres
Advisor
                   : This
                            synchronization
Purpose
                                               primitive
                 threshold values corresponding to
compares
           the
                                                          the
processes P1 through P3, beginning from P1, to FOURTHEVC
and returns the name of the first one which is equal to or
greater than that value. If non of the processes meet this
conditions then P4 is returned.
```

D. THRESH.PLI

SANTANIA TERRETARIA DIDICATA PARAMENTARIA PERCENSIANA PERCENSIANA

```
Prog
                   : THRESH.PLI
    Name
Date
                   : December 83
Written by
                   : M. Kadri Ozyurt
For
                   : Thesis
                   : Professor Kodres
Advisor
                                              primitive
                   : This
Purpose
                            synchronization
receives a pointer to the calling process and increments
the corresponding threshold value by an assigned amount
THRESH: PROC(1);
```

DCL THRESHOLD(0:2)FIXED BIN(15) EXTERNAL,
i FIXED BIN(7);

IF (1=1) THEN THRESHOLD(0)=THRESHOLD(0) + 1;
IF (1=2) THEN THRESHOLD(1)=THRESHOLD(1) + 4;
IF (1=3) THEN THRESHOLD(2)=THRESHOLD(2)+4;
RETURN;
END THRESH;

```
E. P1.PLI
/*
Prog Name
                       : P1.PLI
                       : December 83
Date
Written by
                       : M. Kadri Ozyurt
For
                       : Thesis
Advisor
                       : Professor Kodres
Purpose
                       : This process is basically
infinitive loop. Once entered, it first call AWAIT.PLI
to see if FOURTHEVC is equal to or greater than its
threshold value. If it is, then the control proceeds to
call TACTICAL.PLI. The last call in the loop is to THRESH.PLI to increment its threshold value. In the next iteration, the control will not come back since its
threshold value is greater than FORTHEVC.
```

```
P1: PROCEDURE;

DCL AWAIT ENTRY (FIXED BIN(7)),
    THRESH ENTRY (FIXED BIN(7)),
    A FIXED BIN(7),
    TACTICAL ENTRY;

A=1;
DO WHILE ('1'B);
    CALL AWAIT (1);
    CALL TACTICAL;
    CALL THRESH (1);
END /*DO*/;
END P1;
```

```
F. P2.PLI
/*
Prog Name
                     : P2.PLI
Date
                     : December 83
Written by
                     : M. Kadri Ozyurt
For
                     : Thesis
Advisor
                     : Professor Kodres
                     : The purpose of this process is idecti-
Purpose
cal to that of P1.PLI with the exception that the second
call is to DISPLAY.PLI
P2: PROCEDURE;
  DCL AWAIT ENTRY (FIXED BIN(7)),
THRESH ENTRY (FIXED BIN(7)),
A FIXED BIN(7),
      DISPLAY ENTRY;
  A=2;
  DO WHILE ('1'B);
           CALL AWAIT (2);
           CALL DISPLAY;
           CALL THRESH (2);
  END /*DO*/;
END P2;
```

```
G. P3.PLI
/*
Prog Name
                     : P3.PLI
Date
                     : December 83
                     : M. Kadri Ozyurt
Written by
For
                     : Thesis
Advisor
                     : Professor Kodres
                     : The purpose of this process is the
Purpose
same as P1.PLI with the exception that the second call is
to STATUS.PLI
P3: PROCEDURE;
  DCL AWAIT ENTRY (FIXED BIN(7)),
THRESH ENTRY (FIXED BIN(7)),
      A FIXED BIN(7).
      STATUS ENTRY;
  A=3;
  DO WHILE ('1'B);
           CALL AWAIT (3);
           CALL STATUS;
           CALL THRESH (3);
  END /*DO*/;
END P3;
```

```
H. P4.FLI
/#
Prog Name
                    : P4.PLI
Date
                    : December 83
Written
        ру
                    : M. Kadri Ozyurt
For
                    : Thesis
Advisor
                    : Professor Kodres
Purpose
                    : This process is an infinitive loop
in which there is only one call to IDLE.PLI repeaditively
until an interrupt comes along.
P4: PROCEDURE;
  DCL AWAIT ENTRY (FIXED BIN(7)),
      THRESH ENTRY (FIXED BIN(7)).
      A FIXED BIN(7),
      IDLE ENTRY;
  A=4;
  DO WHILE ('1'B);
          CALL IDLE;
  END /*DO*/;
END P4;
```

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APPENDIX E

MISCELLANEOUS ASSEMBLY ROUTINES

A. KEYBOARD.A86

Prog Name : KEYBOARD.A86 : December 83 ; Da te ;Written by : M. Kadri Ozyurt : Thesis For : Professor Kodres ;Advisor ; Purpose : This program receives a formal iparameter, KEY, reads the status of the serial ;interface chip. If a character has been received from the ;keyboard. it reads this character and places it to the ;formal parameter. If there is not a character available ;it puts a ascii equivalent of zero into the parameter. The reason for that is that zero is not used as a keyboard ; command. The variables used here are defined in the body ; of ARBITER.A86

KEYBOARD:

PUSHF CLI PUSH AX IN AL.PORTST CMP AL . RXRDY KEYBOARD1 JΖ AL, PORTIO IN AL,7FH AND JMP KEYBOARD2 KEYBOARD1: AL.30H MOV KEYBOARD2: BX, [BX] VOM MOV [BX],AL POP AX POPF RET

B. ATOD.A86

;Prog Name : ATOD.A86
;Date : December 83
;Written by : M. Kadri Ozyurt
;For : Thesis
;Advisor : Professor Kodres
;Purpose : This program receives two parameters
;It reads the output of the A/D converter specified by the
;second parameter and places it into the first parameter.
;The variables used here are defined in ARBITER.A86

ATOD:

PUSHF CLI PUSH SI PUSH AX BX PUSH PUSH BX DS PUSH ;BX=.ARGUMENT(1) VOM BX.[BX] MOV . AH,0 AL, [BX] MOV SI,AX MOV VOM AX.SEGCONV MOV DS.AX VOM BX.OFFCONV AL,[BX+SI] :READ A/D PORT VOM POP DS POP BX VOM BX,2[BX] [BX],AL MOV POP BX AX POP POP SI POPF RET

C. RINGBELL.A86

: RINGBELL.A86 ;Prog Name : December 83 ;Date : M. Kadri Ozyurt ;Written by

;For : Thesis

: Professor Kodres ;Advisor

;Purpose : This program sends a bell character to ;the video terminal. The variables used here are defined in ;the body of ARBITER.A86.

RET

RINGBELL:

PUSHF CLI PUSH AX CALL WAIT VOM AL, BEL OUT PORTIO, AL POP AX POPF

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D. WAIT.A86

;Prog Name : WAIT.A86 ; Date : December 83 Written by : M. Kadri Ozyurt

: Thesis

: Professor Kodres ; Advisor

;Purpose : This program program reads the status ;of the serial I/O chip and waits until the transmitter is

; ready to send characters.

WAIT:

PUSH AX

AL, ØDEH WAIT1: IN GET STATUS

AL,1 AND JΖ WAIT1 POP AX

RET

E.SUSPEND.A86

;Prog ;Date Name : SUSPEND.A86 .: December 83 Written by : M. Kadri Ozyurt

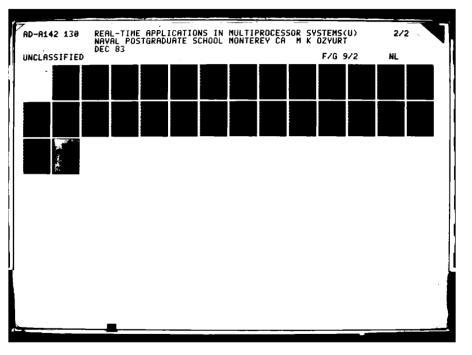
For : Thesis

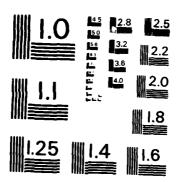
Advisor : Professor Kodres

; Purpose : This program stops the real time clock ; by reseting the interrupt bit of the PSW.

SUSPEND:

CLI RET





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NATIONAL BUREAU OF STANDARDS-1963-A

F. RESUME.A86

¡Prog Name : RESUME.A86
;Date : December 83
;Written by : M. Kadri Ozyurt
;For : Thesis
;Advisor : Professor Kodres
;Purpose : This program starts the real time clock
;by reseting the interrupt bit of the PSW. It then reset
;the counter to zero.

RESUME:

PUSE AX AL, CNTR2 PORTC, AL RESET COUNTER MOV OUT AL, CNTRHI MOV COUNT, AL OUT AL, EOI ; RESET PIC MOV PIC1,AL OUT AX POP STI RET

APPENDIX F

DYNAMIC DEBUGGING MODULE LISTINGS

A. LOCALS AID

/*

Prog Name : LOCALS.AID
Date : December 83
Written by : M. Kadri Ozyurt

For : Thesis

Advisor : Professor Kodres

Purpose : This %include file contains the declarations of the variables used by the dynamic debugging module

*/

SEED TO THE TOTAL SEED TO SEED

DCL
BREAKS(0:9) LABEL,
STOPS(0:9) BIT(1) EXTERNAL,
(CODE1,CODE2,VALUE,H) FIXED BIN(15),
BREAKPT FIXED BIN(15) EXTERNAL,
PUTVARS ENTRY,
REENTRY ENTRY,
BREAKPTS ENTRY,
PROMPTUSER ENTRY (FIXED BIN(7)),
STORESTATUS ENTRY,
TACTICAL ENTRY,
DISPLAY ENTRY,
IDLE ENTRY,
STATUS ENTRY,
CHANGEVA ENTRY (FIXED BIN(15),FIXED BIN(15));

```
B. ERRHAND.AID
```

end /*if*/;

```
/*
Prog
                  : ERRHAND.AID
      Name
                  : December 83
Date
                  : M. Kadri Ozyurt
Written by
                  : Thesis
For
                  : Professor Kodres
Advisor
                  : This %include file contains six
Purpose
different types of PL/I ON condtion bodies. Upon
intercepting any raised error condition is displayed
and the control is transfered to PROMPTUSER with a number
that shows which breakpoint has been past. Then the ON
condition body is exited with a non-local goto statement.
At the exit point the control is transfered to REENTRY
which is the dynamic debugging tool. This call to REENTRY
is protected during the course of normal operation with an
if statement which tests the value of ERRORON.
*/
stops(0)=false;
stops(1)=false;
stops(2)=false;
stops(3)=false;
stops(4)=false;
stops(5)=false;
stops(6)=false;
stops(?)=false;
stops(8)=false;
stops(9)=false;
on error
  begin;
  put list(^Z');
                                 /*clear screen*/
  put skip list('Error #');
/*this statement gets the code of the error condition*/
  code1=oncode();
/*this call prompts the user with the # of breakpoints past
and asks if the user wants to enter the dynamic debugging
environment*/
  call promptuser(code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1<=127 then do;
          put skip list('The program will be abandoned');
          stop;
```

```
end /*error*/;
on fixedoverflow
  begin;
 put list(^Z');
 put skip list('Fixedoverflow #');
  code1=oncode();
 call promptuser (code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1 <= 127 then do;
          put skip list('The program will be abandoned');
          stop;
  end /*if*/;
end /*fixedoverflow*/;
on overflow
  tegin;
 put list(^Z');
  put skip list('Overflow #');
  code1=oncode();
  call promptuser(code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1<=127 then do;
          put skip list('The program will be abandoned');
  end /*if*/;
end /*overflow*/;
on underflow
  begin;
  put list(^Z');
  put skip list('Underflow #');
  code1=oncode();
  call promptuser(code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1<=127 then do;
          put skip list('The program will be abandoned');
          stop;
  end /*if*/;
end /*underflow*/;
on zerodivide
  begin;
```

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```
C. PROMPT.PLI
/*
                       : PROMPT.PLI
Prog Name
Date
                       : December 83
Written by
                       : M. Kadri Ozyurt
For
                       : Thesis
Advisor
                       : Professor Kodres
                       : This routine prompts the user in case
Purpose
of an error interception by ERRHAND. It puts the code of
the error and which breakpoint has been past. Then the
program asks if the user wants to enter the dynamic
debugging environment.
promptuser:proc(code1) external;
/*
  dcl
            %include 'const.inp';
%include 'globals.inp';
            %include 'locals.aid';
  erroron=true;
  put list(code1, 'is detected when the module ');
put skip list('The clock stopped.');
put skip list('The breakpoint #', breakpt, 'has been past.');
put skip list('Would you like to enter interactive debugging',
                        environment (Y/N)?');
  get list(key);
end promptuser;
```

```
/*
                 : REENTRY.PLI
Prog
     Name
                  : December 83
Date
Written by
                  : M. Kadri Ozyurt
For
                  : Thesis
Advisor
                  : Professor Kodres
                  : This routine is the "workhorse" of
Purpose
the dynamic debugging environment. It calls PUTVARS if
the user wants to see the external variables. Then it
calls CHANGEVA if the user wants to change any variable
with in a loop until no changes are wanted. It then transfers
the control to the breakpoint the user desires.
*/
  reentry:proc external;
/*
 dcl
*/
          %include 'const.inp';
          %include 'globals.inp'; %include 'locals.aid';
put skip list (You have entered the interactive debugging
                                               , 'environment.');
   skip list('You will
                             be asked questions about the
put
                                      'control of program flow');
reentry1:
put skip list('Do you want a listing of
                                               all
                                                    variables '
                                                        (Y/N)?');
get list(key);
if (key='Y') ! (key='y') then call putvars();
put skip list('Do you want to change the value of any '
                                               'variable(Y/N)?');
get list (key);
do while ("((key = 'N') ! (key='n')));
  if (key = 'Y') ! (key='y') then do;
  put skip list('Enter the number and the new value',
            (-32768<=value<=+32,767) of the veriable you want to');
       skip list(' change in integers seperated by
                                                        'comma.');
  on error begin;
     put list ('*** bad entry, try again');
     goto reentry2;
```

D. REENTRY.PLI

end /*error*/;

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```
on fixedoverflow begin;
          put list('*** too large, try again');
          go to reentry2;
  end /*fixedoverflow*/;
reentry2:
  put škip list('>');
  get list (code1, value);
  revert error;
  revert fixedoverflow;
  if code1>maxvars then
      put list('invalid variable number');
  else
      call changeva(code1, value);
  /*end if*/
  put skip list('Do you want a listing again (Y/N)?');
  get list(key);
  if (key='Y')!(key='y') then call putvars();
  put skip list ('Do you want to change another
                                                      variable '
                                                         (Y/N)?');
      end /*do*/:
    else
  put list('** bad entry, try again');
    /*end if*/
    get list(key);
end/*do*/;
put skip list('Which breakpoint do you want to transfer the ',
                          control (0 thru 9, foolowed by return)?');
get list(key);
do while ((rank(key)<48) ! (rank(key)>57));
  put list('*** bad entry, try again');
  put skip list('>');
  get list(key);
end/*do*/;
code1=rank(key)-48;
breakpt=code1;
put skip list('Enter the breakpoint you want to stop (0', thru'9) or any non-numeral character if you do not want
                                                    stop(fol. RET');
get list(key);
if ((rank(key)<48) ! (rank(key)>57)) then do;
 put skip list('The program will execute beginning from '
                                           'the breakpoint',code1);
  erroron=false;
      end /*do*/;
else
    do:
  code2=rank(key)-48;
  stops(code2)=true;
  put skip list('The program will execute between the ',
                                'breakpts',code1, and ',code2);
```

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```
end /*if*/;
put skip list('Is that what you want(Y/N)?');
get list(key);
if (key='N') ! (key='n') then do;
  put skip list('Do you want another run(Y/N)?');
  get list (key);
  if (key='Y') ! (key='y') then goto reentry1;
end /*if*/;
```

```
E. PUTVARS.PLI
/*
                : PUTVARS.PLI
     Name
Prog
Date
                : December 83
Written by
                 : M. Kadri Ozyurt
For
                : Thesis
                : Professor Kodres
Advisor
                : This routine puts selected external
Purpose
variables out with PL/I put edit statement.
putvars:proc external;
 %replace
           max_ships by 2,
           true by '1'b.
           false by '0't;
  /*dc1*/
         %include 'globals.inp';
  1=0;
  j=0;
 put skip list('The listing of all common variables is as '
                                                  follows:');
  put skip(2) list('Fixed binary values:');
 put skip edit('(1)seconds=',seconds,'(2)minutes=',minutes,
     (3)hours=',hours,'(4)wake_ptr=',wake_ptr,'(5)I=',i)(r(format1));
 currentproc)(r(format1));
  put skip edit('(14)fourthevc=',fourthevc)(a,f(5));
 put skip(2) list('Fixed decimal values:');
  put skip(2) list('Boolean values:');
 put skip edit('(15)engaged=',engaged,'(16)magnified=',magnified,'(17)fired=',fired,'(18)erroron=',erroron)
               (a,b(1),col(20),a,b(1),col(41),a,b(1),col(58),a,b(1));
  put skip(2) list('Fixed decimal values:');
 put skip edit('(27)vr=',vr,'(28)alpha=',alpha,
  '(29)ax=',ax,'(30)bx=',bx)(r(format2));
put skip edit('(31)cx=',cx,'(32)ay=',ay,'(33)by=',by,'(34)cy=',cy)
```

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(r(format2));

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```
put skip edit(((41)x_at5=',x_at5,'(42)y_at5=',y_at5,'(43)r=',r)(r(format3));
        put skip edit('(47)key=',key)(a(8),a(1));
put skip(2)list('Arrays:');
put skip edit('(48)threshold(0)=',threshold(0),'(49)threshold(1)=',
                     threshold(1), (50)threshold(2)=', threshold(2))
                     (a,f(5),col(26),a,f(5),col(51),a,f(5));
        put skip(2) list('Data structures:');
        put skip list('ship(1):');
        put skip edit('(51)course=',course(1),'(52)speed=',speed(1
        (r(format6));
        put skip list('Ship(2):');
        put skip edit('(63)course=',course(2),'(64)speed=',speed(2),
        '(65)azimuth=',azimuth(2),'(66)range=',range(2))(r(format4));
put skip edit('(67)x=',x(2),'(68)y=',y(2),'(69)x_aim=',x_aim(2),
        '(70)y_aim='.y_aim(2))(r(format5));
put skip edit('(71)count=',count(2),'(72)number=',number(2),
'(73)ptr=',ptr(2),'(74)link_ship=',link_ship(2))
                                                                (r(format6));
        put skip edit('(79)x_wake=',x_wake(ptr(2)),'(80)y_wake='
                y_wake(ptr(2)), (81)link_wake=',link_wake(ptr(2)))
        (a.f(7,1),col(20),a.f(7,1),col(39),a,f(1));
put skip edit ('(82)dt=',dt)(a,f(4,2));
format1: format(\bar{a}, f(2), col(16), a, f(2), col(31), a, f(2), col(44), a, f(2), col(59),
                                                          a.f(5));
format2:format(a,f(4,1),col(19),a,f(4,1),col(37),a,f(4,1),col(58),a,
                                                                  £(4,1));
format3:format(a,f(4,1),col(24),a,f(4,1),col(47),a,f(4,1));
format4:format(a,f(4,1),col(19),a,f(4,1),col(35),a,f(4,1),col(54),a,
format5:format(a,f(7,1),col(19),a,f(7,1),col(35),a,f(7,1),col(54),a,f(7,1));
                                                                     f(7,1));
format6:format(a,f(1),col(19),a,f(2),col(35),a,f(2),col(54),a,f(2));
end putvars;
```

CONTROL CONTROL OF THE CONTROL OF TH

```
F. CHANGEVA.PLI
/*
                 : CHANGEVA.PLI
Proz Name
                 : December 83
Date
                 : M. Kadri Ozyurt
Written by
For
                 : Thesis
Advisor
                 : Professor Kodres
                 : This routine changes a selected
Purpose
external variable specified by the parameter passed
changeva:proc(code1, value) external;
  dcl
    (code1, value) fixed bin(15);
    %include 'const.inp';
    %include 'globals.inp';
  if code1=1 then seconds=binary(value,7);
  if code1=2 then minutes=binary(value,7);
  if code1=3 then hours=binary(value,7);
  if code1=4 then wake_ptr=binary(value,7);
  if code1=5 then i=binary(value,15);
  if code1=6 then t_of=binary(value,7);
  if code1=7 then target=binary(value,7);
  if code1=8 then own=binary(value,7);
  if code1=9 then known=binary(value,7);
  if code1=10 then j=binary(value,15);
  if code1=11 then t=binary(value,7);
  if code1=12 then t_prime=binary(value,7);
  if code1=13 then currentproc=binary(value, ?);
  if code1=14 then fourthevc=binary(value,15);
  if code1=15 then engaged=bit(value,1);
  if code1=16 then magnified=bit(value,1);
  if code1=17 then fired=bit(value,1);
  if code1=18 then erroron=bit(value.1);
  if code1=19 then vx_own=decimal(value,4,1);
  if code1=20 then vy_own=decimal(value,4,1);
  if code1=21 then vx_target=decimal(value,4,1);
  if code1=22 then vy_target=decimal(value,4,1);
  if code1=23 then vx_rel=decimal(value,4,1);
```

if code1=24 then vy_rel=decimal(value,4,1);
if code1=25 then vx_round=decimal(value,4,1);
if code1=26 then vy_round=decimal(value,4,1);

if code1=27 then vr=decimal(value,4,1);

```
if code1=28 then alpha=decimal(value,4,1);
if code1=29 then ax=decimal(value,7,2);
if code1=30 then bx=decimal(value,7,2);
if code1=31 then cx=decimal(value,7,2);
if code1=32 then ay=decimal(value,7.2);
if code1=33 then by=decimal(value,7,2);
if code1=34 then cy=decimal(value,7,2);
if code1=35 then ax_sum=decimal(value,7,2);
if code1=36 then bx_sum=decimal(value,7,2);
if code1=37 then cx_sum=decimal(value,7,2);
if code1=38 then ay_sum=decimal(value,7,2);
if code1=39 then by_sum=decimal(value,7,2);
if code1=40 then cy_sum=decimal(value,7,2);
if code1=41 then x at5=decimal(value.7,2);
if code1=42 then y_at5=decimal(value,7,2);
if code1=43 then r=decimal(value,7,2);
if code1=44 then dx_dt_at5=decimal(value,7,2);
if code1=45 then dy_dt_at5=decimal(value,7,2);
if code1=46 then dr_dt_at5=decimal(value,7,2);
if code1=47 then key=ascii(value);
if code1=48 then threshold(0)=value;
if code1=49 then threshold(1)=value;
if code1=50 then threshold(2)=value;
if code1=51 then course(1)=decimal(value,4,1);
if code1=52 then speed(1)=decimal(value,3,1);
if code1=53 then azimuth(1)=decimal(value,3,0);
if code1=54 then range(1)=decimal(value,5,0);
if code1=55 then x(1)=decimal(value,6,1);
if code1=56 then y(1)=decimal(value.6.1);
if code1=57 then x_aim(1)=decimal(value,6,1);
if code1=58 then y_aim(1)=decimal(value,6,1);
if code1=59 then count(1)=binary(value,7);
if code1=60 then number(1)=binary(value,7);
if code1=61 then ptr(1)=binary(value,7);
if code1=62 then link ship(1)=binary(value,7); if code1=63 then course(2)=decimal(value,4,1);
if code1=64 then speed(2)=decimal(value,3,1);
if code1=65 then azimuth(2)=decimal(value,3,0);
if code1=66 then range(2)=decimal(value,5,0);
if code1=67 then r(2)=decimal(value,6,1);
if code1=68 then y(2)=decimal(value,6,1);
if code1=69 then x_aim(2)=decimal(value,6,1);
if code1=70 then y aim(2)=decimal(value,6,1);
if code1=71 then count(2)=binary(value,7);
if code1=72 then number(2)=binary(value,7);
if code1=73 then ptr(2)=binary(value,7);
if code1=74 then link_ship(2)=binary(value,7);
if code1=75 then az=decimal(value,4,1);
if code1=76 then alt=decimal(value,4,1);
if code1=?? then x_gun=decimal(value,6,1);
```

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```
if code1=78 then y_gun=decimal(value,6,1);
if code1=79 then x_wake(ptr(2))=decimal(value,6,1);
if code1=80 then y_wake(ptr(2))=decimal(value,6,1);
if code1=81 then link_wake(ptr(2))=binary(value,7);
if code1=82 then dt=float(value,7);
```

end changeva;

G. BREAKSØ.AID

/* Prog Name : BREAKSØ.AID : December 83 : M. Kadri Ozyurt Date Written by For : Thesis Advisor : Professor Kodres Purpose : This %include file is one of the ten %include files ,BREAKSØ through BREAKS9, that are used to insert various parts of the programs to be tested. They are protected during the normal operation of the program under test with an if statement. Within the if statement thereis a call to BREAKPTS . */

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```
H. BREAKPTS.PLI
/*
                   : BREAKPTS.PLI
Prog
      Name
                   : December 83
Date
Written by
                   : M. Kadri Ozyurt
                   : Thesis
For
                   : Professor Kodres
Advisor
                   : This routine prompts the user that the
Purpose
breakpoint intended to stop has been reached. Then it asks
if the user wants to transfer the control over the dynamic
debuggig environment. If the answer is positive then it
calls REENTRY where the control stays thereafter.
breakpts:proc external;
  dcl
           stops(0:9) bit(1) external,
           breakpt fixed bin(15) external.
           key char(1) external,
           erroron bit(1) external,
           reentry entry;
  stops(breakpt)='0'b;
  put skip list('**** breakpoint', breakpt, '*****');
  put skip list('The execution halted and clock stopped.'); put skip list('Do you want to enter the interactive debugging',
                    'environment(Y/N)?');
  get list(key);
if (key='Y') ! (key='y') then
```

call reentry();

erroron='Ø'b:

else

/*end if*/
end breakpts;

```
/*
                    : TIMES.AID
Prog Name
                    : December 83
Date
Written by
                    : M. Kadri Ozyurt
                    : Thesis
For
                    : Professor Kodres
Advisor
                    : This %include file is inserted to
Purpose
WAR.FLI to test the execution times of the individual
system routines.
put skip(2) list ('Do you want to measure the execution',
, times of the modules (T/N)?');
get list (key);
do while ((key='Y') ! (key='y'));
  put skip(2) list('Enter the number of iterations you want',
                    '(max 32,767).');
  on error begin;
          put list('*** bad entry, try again.');
          goto times1;
  end /*error*/;
  on fixedoverflow begin;
          put list('** too large, try again.');
          goto times1;
  end /*fixedoverflow*/;
times1:
  put skip(2) list('>');
  get list(h);
  revert error;
  revert fixedoverflow;
  do i=1 to 4;
          put skip(2) list('Ready!! Press any key to start'
          , 'the time check of the module');
if i=1 then put list ('IDLE.');
          else if i=2 then put list('STATUS.');
else if i=3 then put list('TACTICAL.');
          else put list(' DISPLAY.');
          get list(key);
          do j=1 to h while (i=1); call idle; end;
          do j=1 to h while (i=2); call status; end;
do j=1 to h while (i=3); call tactical; end;
          do j=1 to h while (i=4); call display; end;
```

I. TIMES.AID

```
put skip(2)list('The end of the execution .');
put skip(2)list('Enter the time measured in '
                                                  seconds. );
          on error begin;
                  put list('*** bad entry, try again');
                  go to times2;
          end /*error*/;
          on fixedoverflow begin;
                  put list('*** bad entry, try again.');
                  goto times2;
          end /*fixedoverflow*/;
times2:
          put skip(2) list('>');
          get list(j);
          revert error;
          revert fixedoverflow;
          begin;
            dcl duration float;
            duration=float(j)/float(h);
            end;
 end /*do*/;
 put skip(2) list('Do you want another run (Y/N)?');
 get list(key);
end /*do*/;
```

APPENDIX G

A SAMPLE SUPROUTINE TESTING

```
/*
                    : P.PLI
Prog
      Name
                    : December 83
Date
Written by
                    : M. Kadri Ozyurt
For
                    : Thesis
Advisor
                    : Professor Kodres
Purpose
                    : This program is written to test
individual procedures in an interactive manner. At each
bodies are used to intercept any inadvertantly wronger than the terminated either C or
                                     The PL/I
                                                    condition
iteration new values are asked.
                                               ON
entries. The endless loop can be terminated either
   from the terminal. In this particular
                                                 example the
procedure DRAW inside the body of DISPLAY.PLI is tested by
making it external for the test purposes.
p:proc options(main);
dcl
  (u,v)(0:10) fixed bin(15),
  (i,x,y) fixed bin(15),
  rub char(1) external,
  d fixed bin(7),
  draw entry ((0:10) fixed bin(15), (0:10) fixed bin(15)
                   .fixed bin(7));
on error begin;
    put skip list('*** bad value, try again');
    goto reentry;
end;
on fixedoverflow begin;
    put skip list('*** too large, try again');
    goto reentry;
end;
```

```
reentry:
rub=ascii(127);
do while('1'b);
   put list('^/^ X');
   put skip list('enter x and y');
   put skip list('>');
   get list(x,y);
   put list('2');
   call cen(x x y x);
      call gen(x,y,u,v);
       d=1:
      call draw(u,v,d);
call delay;
      d=0:
       call draw(u.v.d);
end;
delay:proc;
   dcl (i,j) fixed bin(15);
   do i=1 to 30000;
                do j=1 to 2; end;
end;
end;
gen:proc(x,y,u,v);
   dcl
                 (u,v)(0:10) fixed bin(15), (x,y) fixed bin(15);
                             v(0)=y;
v(1)=y+8;
u(0)=x+8;
u(1)=x;
                             v(2)=y;
u(2)=x-8;
                             v(3)=y-8;
u(3)=x;
                              v(4)=y;
u(4)=x+8;
u(5)=-1;
                              v(5)=-1;
end;
end p;
```

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的时间*是是各种的*是是是有的,我们是是是是是是一种,我们是是是是是是是一种,我们也是是是是是是是是是是是是是是是是是是,我们也是是是是一种,我们也是是是一种的人,

APPENDIX H A SAMPLE PROGRAM TESTING

/* Prog Name : IDLE.PLI Date : December 83 : M. Kadri Ozyurt Written by : Thesis For : Professor Kodres Advisor Purpose : This is the testing version of the procedure IDLE.PLI under the dynamic debugging module. After the correct result from the test had been taken, the final version of the procedure was made simply removing the segment of code in between the comment lines. In order to test the program, an interactive main procedure as in Appendix G was written.

IDLE: PROCEDURE EXTERNAL;

THE PROPERTY OF THE PROPERTY O

```
/*
DCL

*/

ZINCLUDE 'CONST.INP';

ZINCLUDE 'GLOBALS.INP';

/******** DEBUG AID ********/

ZINCLUDE 'LOCALS.AID';

ZINCLUDE 'ERRHAND.AID';

ZINCLUDE 'BREAKSØ.AID';

/*********************

DO D=Ø TO 5;

CALL ATOD (D,ARG(D));

END /*DO*/;
```

```
/*at this point the A/D output values are fixed bin(7)
values.
         The following sequence converts those to fixed
decimal values*/
/***** DEBUG AID *****/
%INCLUDE 'BREAKS1.AID';
/***** END AID ******/
      COURSE(OWN) = ARG(0);
      SPEED(OWN)=ARG(4);
      COURSE(KNOWN) = ARG(2);
      SPEED(KNOWN)=ARG(3);
      AZ=ARG(1);
      ALT=ARG(5);
/**** DEBUG AID *****/
     %INCLUDE 'BREAKS2.AID';
/***** END AID ******/
/*the following sequence converts A/D values to real time
values by using appropriate proportionality constants*/
      COURSE(OWN) = COURSE(OWN) * K;
      COURSE(KNOWN) = COURSE(KNOWN) * K;
      AZ = AZ * K;
      IF COURSE(OWN)<0.0 THEN
          COURSE(OWN) = COURSE(OWN) + TWO_PI;
      IF COURSE (KNOWN) < 0.0 THEN
          COURSE(KNOWN) = COURSE(KNOWN) + TWO PI;
      IF AZ<0.0 THEN
          AZ = AZ + TWO PI;
       IF ALT>90.0 THEN
          ALT = 90.0;
/***** DEBUG AID *****/
     %INCLUDE 'BREAKS3.AID';
/***** END AID ******/
          SPEED(OWN) = SPEED(OWN)/L;
          SPEED(KNOWN) = SPEED(KNOWN) / L;
/*ownship speed computations*/
      VX OWN = SPEED(OWN) * SIND(COURSE(OWN));
      VY OWN = SPEED(OWN) * COSD(COURSE(OWN));
/***** DEBUG AID *****/
     %INCLUDE 'BREAKS4.AID';
/***** END AID ******/
```

```
/*when not have fired, the following makes the tallistic
computations*/
     IF
          FIRED THEN
        BEGIN;
         T_OF = 2.0 * VM * SIND(ALT) / G;
/**** DEBUG AID *****/
    %INCLUDE 'BREAKS5.AID';
/***** END AID ******/
         VR = VM * COSD(ALT);
         R = VR * T OF;
/***** DEBUG AID *****/
    %INCLUDE 'BREAKS6.AID';
X_{AIM}(OWN) = R * SIND(AZ);
         TAIM(OWN) = R * COSD(AZ);
         X = 0.0;
         Y = GUN = 0.0;
/**** DEBUG AID *****/
    %INCLUDE 'BREAKS7.AID';
/****** END AID ******/
         VX ROUND = VR * SIND(AZ);
         VYROUND = VR + COSD(AZ);
/**** DEBUG AID *****/
    %INCLUDE 'BREAKS8.AID';
/***** END AID ******/
         T = T_OF;
     END /*IF*/;
/***** DEBUG AID *****/
    %INCLUDE 'BREAKS9.AID';
/***** END AID ******/
END IDLE;
```

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